

Quadratic Transformations: Convexity vs Nonconvexity

B. Polyak and P. Scherbakov

Weekly seminar, Lab 7

Institute for Control Science, Moscow

February 17, 2015

Problem

Have: Quadratic mapping $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$ of the form

$$f(x) = (f_1(x), \dots, f_m(x))^\top, \quad f_i(x) = (A_i x, x) + 2(b_i, x) + c_i, \quad i = 1, \dots, m \leq n$$

Consider

$$E = \{f(x) : x \in \mathbb{R}^n\}$$

or

$$F = \{f(x) : x \in \mathbb{R}^n, \|x\| \leq 1\}$$

Is F convex or not?

Note: The A_i s are not assumed to be sign-(semi)definite!

Applications

- General quadratic programming:

$$\min f_0(x)$$

$$\text{s.t.} \quad f_i(x) \leq 0, \ i \in I, \quad f_i(x) = 0, \ i \in J$$

- Integer optimization

$$x_i = \{-1, +1\} \iff x_i^2 = 1$$

- Convex relaxation: When is it tight?
- S -theorem (absolute stability): When do the two quadratic inequalities imply the third one?
- Robustness analysis (so-called real μ)
- Power systems: Electric power is a quadratic function of the currents or voltages
- Quantum systems

Convexity vs Nonconvexity

- Simplest example:

$$\min(Ax, x) \quad \text{over} \quad \|x\| = 1$$

This problem is **nonconvex**! However the solution is explicit:

$$x^* = e_1,$$

where e_1 is the eigenvector associated with the minimal eigenvalue of A

- Titles of papers:
 - *Hidden convexity in some nonconvex quadratically constrained quadratic programming* [Ben-Tal, Teboulle, 1996]
 - *Permanently going back and forth between the “quadratic world” and the “convexity world” in optimization* [J.-B. Hiriart-Urruty, M. Tork, 2002]
- When the images of quadratic maps are convex?

Simple Illustrations

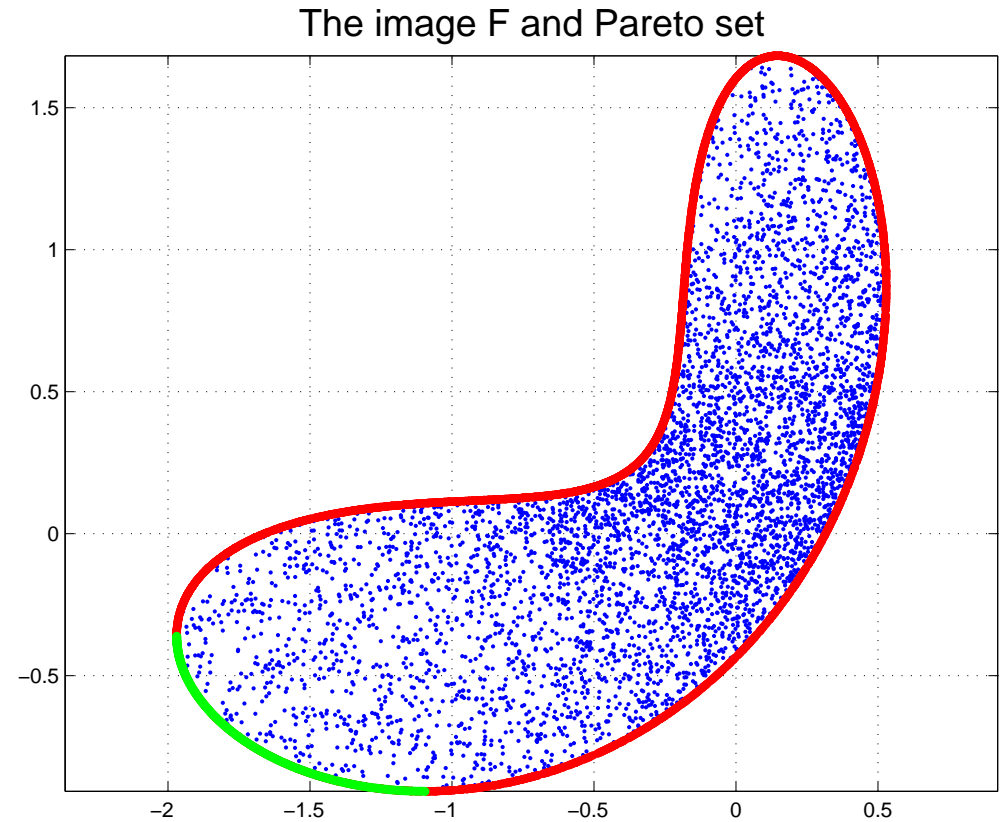
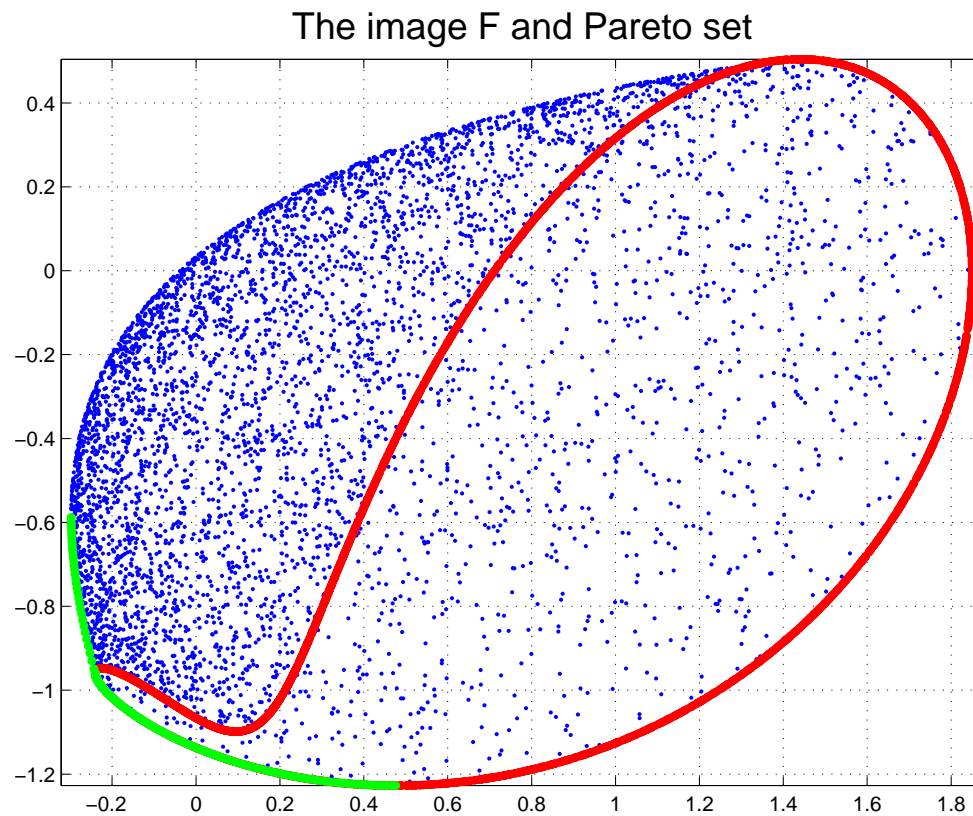


Figure 1: $n = m = 2$: Surface (red) and interior (blue)

Known Facts

- $m = 2, b_i = 0 \implies E$ is convex [Dines, 1941]
- $m = 2, b_i = 0, n \geq 3 \implies F$ is convex [Brickman, 1961]
- $m = 3, b_i = 0, n \geq 3, \sum c_i A_i \succ 0 \implies E$ is convex [Polyak, 1998]
- m is arbitrary, $b_i = 0, A_i$ commute $\implies F$ is convex [Fradkov, 1973]
- m is arbitrary, $b_i \neq 0 \implies F_\varepsilon$ is convex for $\varepsilon > 0$ small enough [Polyak, 2001]

Convex Hull (i)

The idea of convex relaxations for quadratic problems goes back to [Shor, 1986]; see also [Nesterov 1998], [Zhang 2000], [Beck and Teboulle, 2005].

Recent survey:

Luo, Ma, So, Ye, Zhang, *Semidefinite relaxation of quadratic optimization problems*, IEEE Sig. Proc. Magazine, 2010.

Two typical results:

Lemma 1. For $b_i = 0$ have

$$\text{Conv}(F) = \{\mathcal{A}(X) : X \succeq 0, \text{Tr} X \leq 1\},$$

where $X = X^\top \in \mathbb{R}^{n \times n}$, $\mathcal{A}(X) = (\langle A_1, X \rangle, \dots, \langle A_m, X \rangle)^\top$, and $\langle A, X \rangle = \text{Tr} AX$.

Convex Hull (ii)

Lemma 2. *In the general case ($b_1 \neq 0$) have*

$$\text{Conv}(F) = \{\mathcal{H}(X): X \succcurlyeq 0, \text{Tr}X \leq 1, X_{n+1,n+1} = 1\}$$

where $X = X^\top \in \mathbb{R}^{(n+1) \times (n+1)}$, $\mathcal{H}(X) = (\langle H_1, X \rangle, \dots, \langle H_m, X \rangle)^\top$,

$$\text{and } H_i = \begin{bmatrix} A_i & b_i \\ b_i^\top & c_i \end{bmatrix}.$$

Idea of proof: $(A_i x, x) = \langle A_i, x x^\top \rangle = \langle A_i, X \rangle$, $X \succcurlyeq 0$, $\text{rank} X = 1$, $\text{Tr} X = \|x\|^2$.

For $z = (x; t) \in \mathbb{R}^{n+1}$ have $(H_i z, z) = (A_i x, x) + 2(b_i, x)t + c_i t^2 = f_i(x)$ if $t = 1$.

How Close are F and $\text{Conv}(F)$?

- the 0.87856 result for MaxCut [Goemans and Williamson, 1995]
- entropy distance [Barvinok, 2014]

However in general F and $\text{Conv}(F)$ can differ a lot.

Example:

$$n = m = 3, \quad f_1(x) = x_1x_2, \quad f_2(x) = x_1x_3, \quad f_3(x) = x_2x_3$$

FIGURE (four-dent image)

Nonconvexity Certificate N1

Assume

- $b_i = 0$
- the LMI $\sum c_i A_i \succcurlyeq 0$ has no nonzero solution c
- there exists $h \notin E$.

Then E is nonconvex.

Explanation: E is a cone; if it is convex it is either coincides with the whole space
or lies in a half-space

Example: $\text{Tr} A_i = 0$ for all i as above

Similar result holds for $b_i \neq 0$

Convexity Certificate C1

Assume that for all $c \neq 0$, the minimal eigenvalue of the matrix $\sum c_i A_i$ is simple.

Then F is convex. If the eigenvalue is negative, then F is strictly convex.

Indeed, for $G = \text{Conv}(F)$ have

$$\min_{g \in G}(c, g) = \min_{X \succeq 0, \text{Tr} X \leq 1} \langle A, X \rangle;$$

the unique minimizer is either $X = e_1 e_1^\top$ (if $\lambda_1 < 0$) or $X = 0$ (if $\lambda_1 \geq 0$).

Here, $\lambda_1 = \lambda_{\min}(A)$, $A = \sum c_i A_i$, and e_1 is the corresponding eigenvector.

Hence, any supporting hyperplane has a unique intersection point both with F and its convex hull.

Corollary

If the A_i s are matrices with positive entries, the “positive part” of F is convex.

Extension [Jetakumar et al., 2009]). The matrix A is referred to as a z -matrix if $a_{ij} \leq 0, i \neq j$.

If all A_i s are z -matrices, the Pareto set in the multi-objective minimization problem with maps

$$f_i(x) = (A_i x, x)$$

is convex.

Local Convexity

Theorem 1. *Let $b_i = 0$ and the eigensystem of the matrix $A^c = \sum c_i A_i \succcurlyeq 0$ be*

$\lambda_1 < \lambda_2 \leq \dots \leq \lambda_n$, $\lambda_1 < 0$, and e_1, \dots, e_n . Denote $x^c = e_1$, $f^c = f(x^c)$.

Then f^c is the boundary point of F , and for all d such that

$$\|d - c\|_\infty \leq \varepsilon < \frac{\min\{-\lambda_1, \lambda_2 - \lambda_1\}}{2\kappa L};$$

f^d is also a boundary point of F .

Here, $L = \sum \|A_i\|$, $\kappa = \|T\| \|T^\top\|$, where T is the matrix with columns e_i , $T^{-1} = T^\top$, and the matrix norm is $\|M\| = \max_i \sum_j |m_{ij}|$, while the vector f^d is defined similarly to f^c .

In principle, this can provide us with a global convexity certificate: Sample points c^1, \dots, c^N (e.g., randomly on the unit sphere $\|c\| = 1$). If, for each c^i , Theorem 1 holds, and the corresponding small boxes cover the unit sphere, then F is convex.

Nonconvexity Certificate N2

Theorem 2. Let $m \geq 3$, $n \geq 3$, $b_i = 0$, and let for some c_1, \dots, c_m , the matrix $A^c = \sum c_i A_i$ has eigenvalues $\lambda_1 = \lambda_2 < \lambda_3 \leq \dots \leq \lambda_n$, $\lambda_1 < 0$,

and eigenvectors e_1, \dots, e_n .

If $f^1 = f(e_1) \neq f^2 = f(e_2)$, $f^0 = ((A_1 e_1, e_2), \dots, (A_m e_1, e_2))^\top \neq 0$,

then F is nonconvex.

Proof: $\text{Arg min}_{f \in F}(c, f) = f^\alpha$, where $f^\alpha = f(x^\alpha)$; $x^\alpha = e_1 \alpha_1 + e_2 \alpha_2$; $\alpha_1^2 + \alpha_2^2 = 1$.

The set of such points f^α is a 2D ellipse, which is nondegenerate due to the assumptions $f^1 \neq f^2$, $f^0 \neq 0$.

Hence, the intersection of F and the supporting hyperplane $(c, f) = \lambda_1$ is nonconvex.

Example

$$n = m = 3, \quad A_1 = \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}, \quad A_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 2 \end{pmatrix}, \quad A_3 = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix},$$
$$b_1 = b_2 = b_3 = 0$$

For $c = (0, -1, -1)^\top$ we have

$$\lambda_1 = \lambda_2 = -3, \quad \lambda_3 = -2, \quad e_1 = (1, 0, 0)^\top, \quad e_2 = (0, 0, 1)^\top$$

$$f^1 = (0, 1, 2)^\top, \quad f^2 = (0, 2, 1)^\top, \quad f^0 = (1, 0, 0)^\top,$$

hence, Theorem 2 holds and F is nonconvex.

FIGURE (concave ice-cream cone)

How to Find Such c ?

Choose $a \neq 0$ and solve SDP in the scalar variables c, μ :

$$\begin{aligned} & \max \mu \\ \text{s.t.} \quad & \sum c_i A_i \succcurlyeq \mu I \\ & (c, a) = 1 \end{aligned}$$

It is likely that μ is the minimal eigenvalue of $\sum c_i A_i$ with multiplicity 2.

Then check other conditions of Theorem 2 for this c to assure the nonconvexity of F

Otherwise try another a .

Another Approach

Sample points on the boundary of $G = \text{Conv}(F) = \{\mathcal{A}(X) : X \succcurlyeq 0, \text{Tr} X \leq 1\}$:

- If $\sum c_i A_i \succcurlyeq 0$ is infeasible, have $G = \text{Conv}(F) = \{\mathcal{A}(X) : X \succcurlyeq 0, \text{Tr} X = 1\}$.

Take $X^0 = (1/n)I$, pick a random matrix “direction” H , and compute

$$X^1 = X^0 + \gamma_1 \left(H - \frac{1}{n} \text{Tr} H \cdot I \right), \quad \gamma_1 = \max \left\{ \gamma : X^0 + \gamma \left(H - \frac{1}{n} \text{Tr} H \cdot I \right) \succcurlyeq 0 \right\}$$

Then X^1 is a random point on the boundary of G .

Compute the normal vector $c \neq 0$ to G at X^1 by solving the LMI

$$\sum c_i A_i \succcurlyeq \mu I, \quad \left\langle \sum c_i A_i - \mu I, X^1 \right\rangle = 0$$

- If $\sum c_i A_i \succcurlyeq 0$ is feasible, set $X^0 = (1/2n)I$ and proceed similarly without scaling $\text{Tr} X^1 = 1$.

If X^1 is not rank one, then it is likely to sit at the “flat” part of the boundary of G , and we are within the framework of Theorem 2.

Otherwise, keep on generating the points X^k similarly.

Different Problem Formulation

The goal is to approximate F or its boundary point-wise.

Brute force method: Monte Carlo. Sample points x^k in the unit ball $\|x\| \leq 1$ and calculate $f^k = f(x^k)$; adopt as the desired approximation.

OK for small n but very poor for n large.

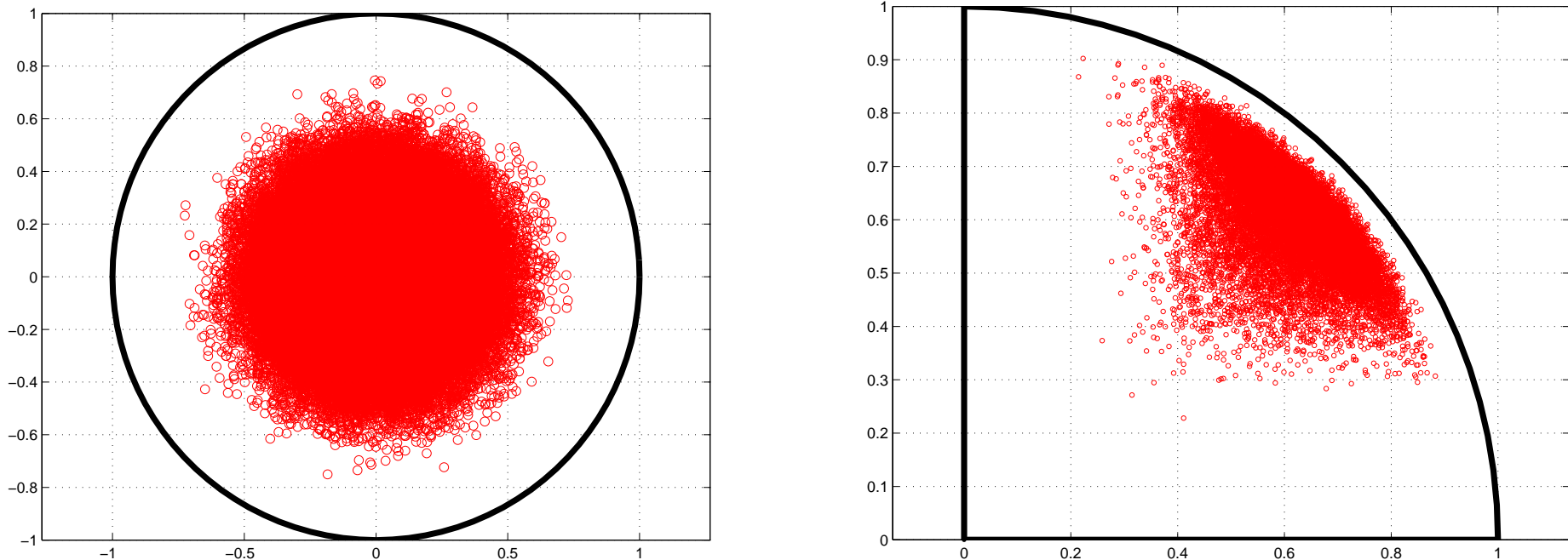


Figure 2: $n = 25$: Linear (left) and homogenous diagonal quadratic (right) mappings.
True boundary and $N = 100,000$ MC samples

A Necessary Condition for a Boundary Point

Lemma 3 (The homogeneous case). If f^* is a boundary point of F , then there exists a $c \neq 0$ such that $f^* = f(e_i)$, e_i being an eigenvector of $A = \sum c_i A_i$ corresponding to a positive or zero eigenvalue of A .

Similar result holds for the nonhomogeneous case.

Hence, for m small, we can sample $c^k \in \mathbb{R}^m$, $\|c^k\| = 1$,

compute A^k s and their eigenvectors and

the points f^k .

However this is not efficient for $m > 2$.

2D image

Special case $m = 2$ is worth analyzing.

Of course only nonhomogenous maps are of interest (in view of Brickman theorem)

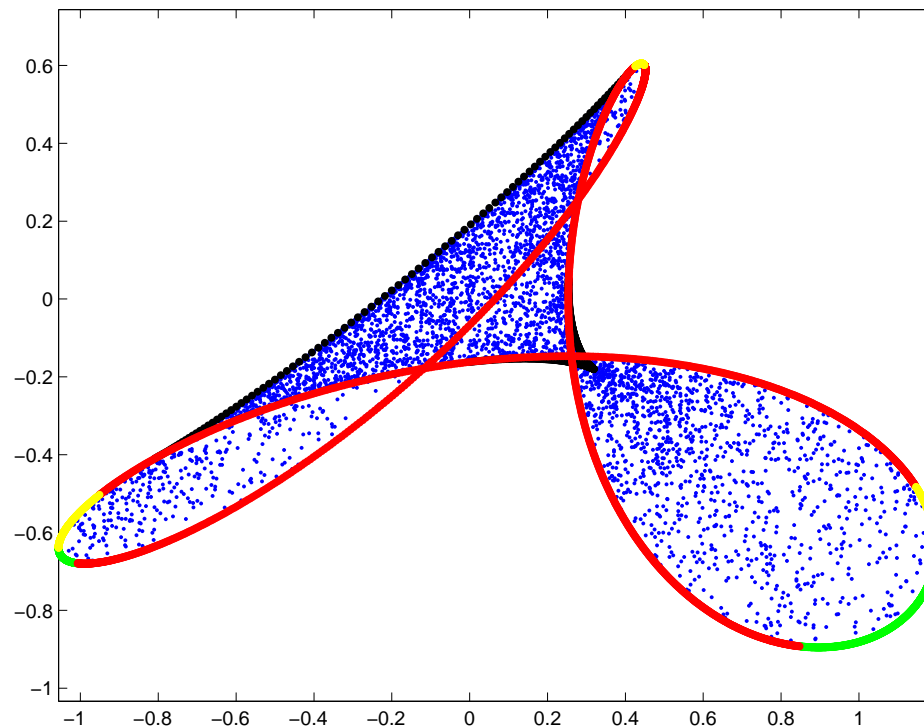


Figure 3: $n = 3, m = 2$

Future Work