

# Geometrical Intuition and Failure of Monte-Carlo in Multidimensional Optimization

*Boris Polyak*

with Pavel Scherbakov, Elena Gryazina

Institute for Control Science, Moscow

Seminar MIPT

March, 2017

# Content

Blend of 3 topics

- Geometry of  $n$  dimensions
- Monte-Carlo method
- Optimization: MC and other grids

## $\mathbb{R}^n$ contradicts our intuition

- Volume of  $n$  dimensional unit ball tends to zero

$$V_n(B_r) = v_n r^n, r \text{ radius of the ball } B_r = \{x : \|x\| \leq r\},$$

$$v_n = \frac{2\pi}{n} v_{n-2}, v_1 = 2, v_2 = \pi, \max_n v_n = v_5 = 5.26$$

- Unit cube  $C = \{x : |x_i| \leq 0.5\}$  is mostly outside unit ball:  $\max_{x \in C} \|x\| = \sqrt{n}/2$
- The ball inscribed in the cube is a small part of its volume (ratio is  $v_n 2^{-n}$ ).
- Volume of a ball is concentrated near the surface:

$$V_n(B_{1-1/n})/V_n(B_1) = (1 - \frac{1}{n})^n \simeq \frac{1}{e}$$

- Center of gravity  $x^*$  of a cone  $K$  lies near the bottom:

$$K = \{x : 1 \geq x_1 \geq \alpha \sqrt{x_2^2 + \dots + x_n^2}\}, x_1^* = 1 - \frac{1}{n+1}$$

- Volume of spherical cap is small:  $H = \{x : x_1 \geq h < 1, \|x\| \leq 1\}$ .  $V_3(H)$  was calculated by Archimedes, for large  $n$ ,  $h = \frac{\varepsilon}{\sqrt{n-1}}, \frac{V_n(H)}{V_n(B_1)} \leq f(\varepsilon)$ .

# The $\mathbb{R}^n$ geometry

John Hopcroft and Ravindran Kannan,

*“Foundations of Data Science, 2014,”* Chapter 2, High-Dimensional Space

# Monte Carlo: Beginning

J. von Neumann, E. Teller, S. Ulam and N. Metropolis (1949).

## First applications

- Simulation (physical processes of particle motion and binary collisions).
- Multidimensional integration.
- Computational mathematics.

## Recent applications

- Convex and global optimization
- Robustness analysis in control and optimization.
- Sampling in complex domains.
- Data mining, Learning, Signal processing, Image processing

# Goal: Sampling in complicated domains

Possible approaches:

- Special regions (box, simplex,  $l_p$ -ball, positive definite matrices cone)
- Rejection
- Markov chain Monte Carlo (MCMC): random walk in the domain

# 1D Sampling

**Long ago:** Tables of random numbers, physical devices (analogues of roulette or coin).

**Present:** there exist highly effective pseudo-random generators, fast and with good statistical properties.

Example: routine `rand` in Matlab generates points uniformly distributed in  $[0, 1]$ .

How to generate random variables with a given cdf  $F(x)$ ?

**Solution:**  $x = \text{rand}$ ,  $y = F^{-1}(x)$ .

Gaussian  $\mathcal{N}(0, 1)$ : use routine `randn`; alternatively

$$x_1 = \text{rand}, x_2 = \text{rand},$$

$$y_1 = \sqrt{-2 \log x_1} \cos(2\pi x_2),$$

$$y_2 = \sqrt{-2 \log x_1} \sin(2\pi x_2),$$

Then  $y \sim \mathcal{N}(0, I_2)$ .

# nD Sampling

1.  $\text{rand}(n, 1)$  generates the uniform distribution over the box  $B = [0, 1]^n \in \mathbb{R}^n$ .
2.  $\text{randn}(n, 1) \sim N(0, I_n)$ .

**Exercises:** How to generate samples uniformly distributed in:

1) sphere; 2) ball; 3) simplex; 4)  $l_p$ -ball; 5) ellipsoid; 6) matrix balls

**General problem:** given a bounded set  $Q \in \mathbb{R}^n$ , how to generate points uniformly distributed in  $Q$ ? Typical example:  $Q$  is a polytope.

**Rejection:** Take simple  $G \supseteq Q$  (e.g. a box), generate points uniformly in  $G$ , reject those which are not in  $Q$ .

Bad:  $Q = B_1$ ,  $G = [-1, 1]^n$ . Then  $V_n(G)/V_n(Q) = v_n 2^{-n} = 4 \cdot 10^7$  for  $n = 20$ , that is we should generate  $10^8$  points to get 1–2 points in  $Q$ . For polytopes this ratio can be much larger.



# Markov Chain Monte Carlo (MCMC)

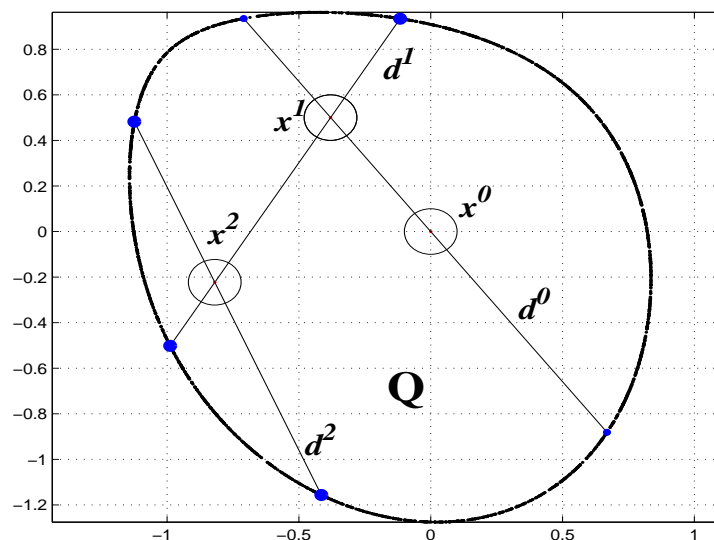
Idea: Apply random walks (vs independent generation of points as in MC).

**P.Diaconis**, The Markov Chain Monte Carlo Revolution, Bull. AMS, Vol. 46, No. 9, pp. 179–205 (2009).

A lot of applications in applied math.

# Hit-and-Run algorithm (sampling in $Q \subset \mathbb{R}^n$ )

This is random walk in  $Q$ . Turchin (1971), Smith (1984)



1. Initial point  $x^0 \in Q$ .
2.  $d = s/\|s\|$ ,  $s = \text{randn}(n, 1)$  — random direction uniform on the unit sphere
3. **Boundary oracle:**  $L = \{t \in \mathbb{R} : x^0 + td \in Q\}$
4. Next point  $x^1 = x^0 + t_1 d$ ,  $t_1$  is uniform random in  $L$ .
5.  $x^0$  is replaced by  $x^1$ , go to Step 2.

## Hit-and-Run: properties

**Theorem:** The points  $x^k$  are asymptotically uniform in  $Q$ , i.e., the probability to reach a subset  $A \subset Q$  can be estimated as:  $|P_k(A) - P(A)| \leq cq^k$ ,

where  $P(A) = \text{Vol}(A)/\text{Vol}(Q)$ ,  $P_k(A) = (\text{number of times } x \text{ visits } xA)/k$

where  $q < 1$  does not depend on the initial point  $x^0$  (but it depends on the dimension and the geometry of the set).

Idea: for any sets  $A \subset Q$ ,  $B \subset Q$  with nonzero volume the probabilities of transition are equal and positive.

# Advantages of HR

- Very simple.
- $Q$  should not be neither convex nor connected (boundary oracle may consist of several intervals), the only assumption is that  $Q$  is a closure of an open set.
- Points on the boundary are also generated.
- Boundary oracle is available for numerous sets.

Extensions and different sampling algorithms ([Shake-and-Bake](#), [Billiard walk](#)).

Polyak B., Gryazina E., “Random sampling: Billiard Walk algorithm”, European Journal of Operational Research, 2014.

# Optimization

Two different setups:

- Convex optimization
- Global optimization

# Convex optimization

Advanced theory, efficient methods + validation, available software for various dimensions. No need for Monte-Carlo.

$\min f(x), x \in Q \subset \mathbb{R}^n, f, Q - \text{convex}.$

“Best method”: centers of gravity method.

Estimate  $\min_{0 \leq i \leq k} (f(x_k) - f^*) \leq cq^k, q \simeq (1 - \frac{1}{ne}).$

Example ( $K$  is the cone), the estimate is tight:

$$\min x_1, x \in K$$

# Global optimization

Numerous heuristic methods (simulated annealing, tabu search, ant colony etc.) Randomization is highly helpful. Negative theoretical results: “Global optimization is unsolvable for  $n \geq 15$ ” Nemirovski, Yudin 1976; Nesterov 2009. Constrained minimization of Lipschitz functions on a cube in  $\mathbb{R}^n$ , a function equals 0 everywhere beyond a cell of a grid; its minimum equals  $-\varepsilon$ . We need  $(\frac{L}{2\varepsilon})^n$  calculations of the function to get into the cell.

Another example:

$$f(x) = \min \left\{ 99 - c^\top x, (c^\top x - 99)/398 \right\}$$

to be minimized over the ball  $B_{100} \subset \mathbb{R}^n$ . It has one local minimum  $x_1 = -100c$ ,  $f_1 = -0.5$ , and one global minimum  $x^* = 100c$ ,  $f^* = -1$ . Any standard method (say, multi-start) misses global optimum with probability larger than  $1 - 10^{-15}$  for  $n = 15$ .

# MC for a ball

Conclusion: for hard problems Monte-Carlo-like methods have no competitors.

Question: how they behave for “good” problems?

Example:  $\min f(x), x \in B, f(x) = (c, x), \|c\| = 1, f^* = -1$

Algorithm:  $x^i, i = 1, \dots, k$  are sample of uniformly distributed points in the unit ball  $B$ ,  
 $f_k = \min_{1 \leq i \leq k} f(x^i)$ .

Estimate How close is  $f_k$  to  $f^*$ ?



## Rigorous result

**Theorem** Given  $p \in ]0, 1[$  and  $\delta \in ]0, 1[$ , the minimal sample size  $N_{\min}$  that guarantees, with probability at least  $p$ , for the empirical maximum of  $f_N$  to be at least a  $\delta$ -accurate estimate of  $f^*$ , is given by

$$N_{\min} = \frac{\ln(1 - p)}{\ln \left[ \frac{1}{2} + \frac{1}{2} I \left( (1 - \delta)^2; \frac{1}{2}, \frac{n+1}{2} \right) \right]},$$

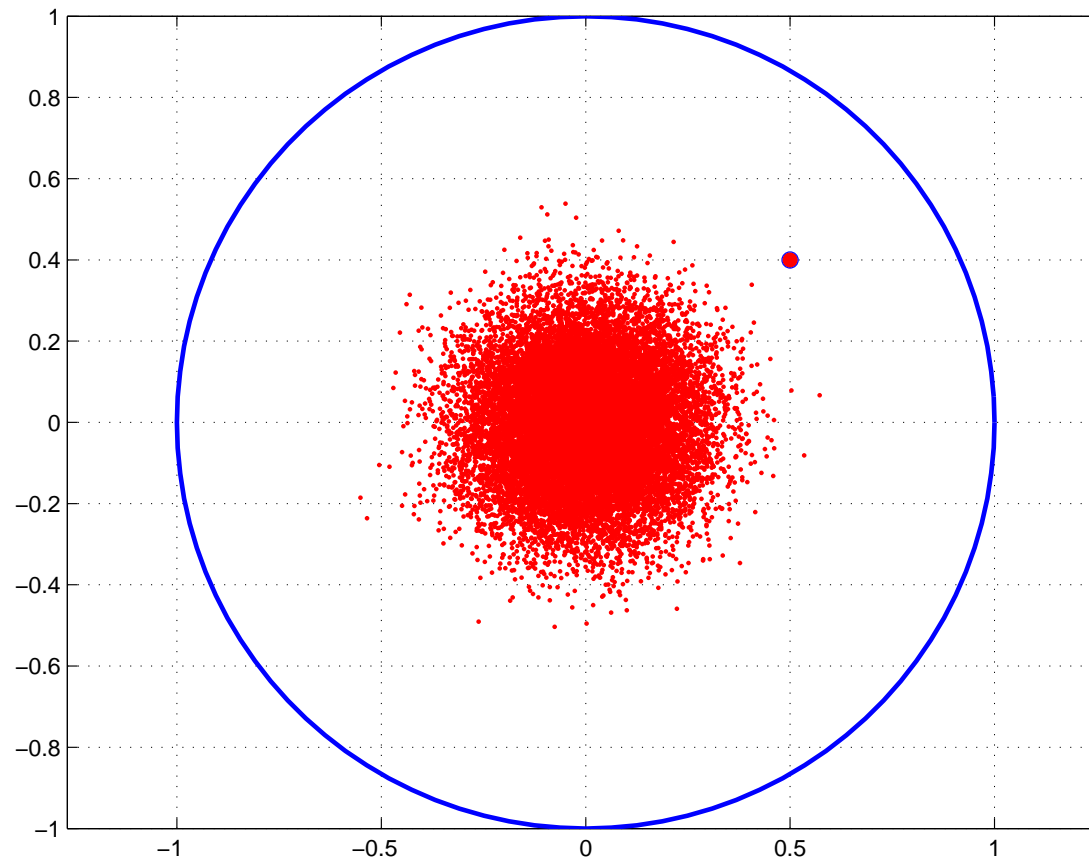
where  $I(x; a, b)$  is the regularized incomplete beta function with parameters  $a$  and  $b$ .

For example for  $n = 10$ ,  $\delta = 0.05$ , and  $p = 0.95$ , it gives  $N_{\min} \approx 8.9 \cdot 10^6$ . There are numerous approximate formulae and other estimates.

# MC for multiobjective optimization

Objective function is  $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$ , construct Pareto front.

Example:  $m = 2$ ,  $f(x)$  is linear,  $B$  is the unit ball,  $n = 50$ ,  $N = 100000$



## MC for a box

$$\min(c, x), x \in Q = [-1, 1]^n, c = (1, \dots, 1)$$

$$N_{\min} = \frac{\ln(1 - p)}{\ln\left(1 - \frac{n^n \delta^n}{2^n n!}\right)}.$$

$n = 10$  ,  $\delta = 0.1$  and  $p = 0.95$ , provide a huge  $N_{\min} \approx 1.12 \cdot 10^{10}$ .

# Deterministic grids

Example  $\max(c, x), x \in Q = [-1, 1]^n, c = (1, \dots, 1)$

**Uniform grid** Uniform grid on  $Q$ , mesh points do not cover the boundary.

**Sobol sequences** ( $LP_\tau$  sequences)

Results for  $N = 10^6$

$n$ ; true max	2	3	4	5	10	15	20
Uniform grid	1.9960	2.9406	3.7576	4.4118	6.0000	7.5000	6.6667
$LP_\tau$	1.9999	2.9792	3.8373	4.6844	7.9330	10.2542	10.9470
Monte Carlo	1.9974	2.9676	3.8731	4.6981	7.8473	10.0796	11.8560

# Literature

B.Polyak, P.Shcherbakov “Why does Monte Carlo Fail to Work Properly in High-Dimensional Optimization Problems?” JOTA, 2017 (available on-line)