

A Tale of Two Models for Random Graphs

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⦿ Random Graph $G(n,p)$

- Component and Branching process
- Emergence of Giant (connected) Component
- Motivation for Poisson cloning
- a. Poisson Cloning Model (PCM)
- b. Cut-Off Line Algorithm (COLA)
- c. New Results via PCM and COLA

⦿ Random Graph $G(n,m)$

- a. Graph Process $G(n,1), G(n,2), \dots, G(n,m), \dots$
- b. Graph Process under constraints
- c. Triangle-free Process and Ramsey Number $R(3,t)$

Complete Graph K_n

- When the edge set E of a graph $G = (V, E)$ is the set of all pairs of distinct vertices, then the graph is called the **complete graph** on V , and denoted by K_V , or simply K_n , where n is the number of all vertices.

Notice that there are

$$\binom{n}{2} := \frac{n(n-1)}{2} \quad \text{edges in } K_n.$$

Random Graph $G(n, p)$

Each of $\binom{n}{2}$ edges is independently
in $G(n, p)$ with probability p ,

$p = 1$: complete graph

$p = 0$: empty graph

- For example, if $n = 4, p = 1/2$, toss a fair coin six times

an instance

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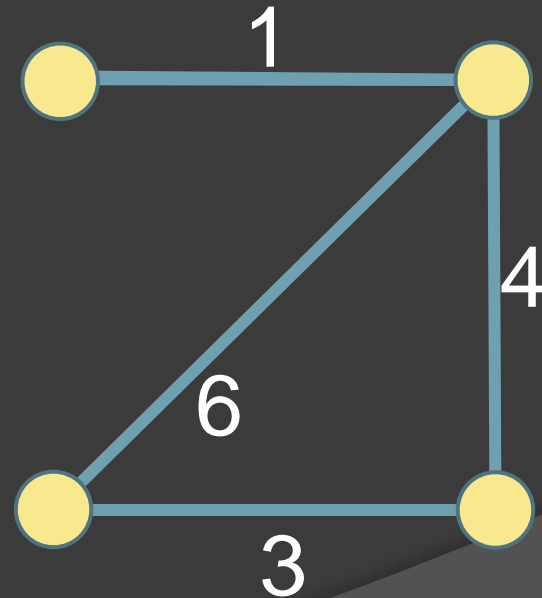
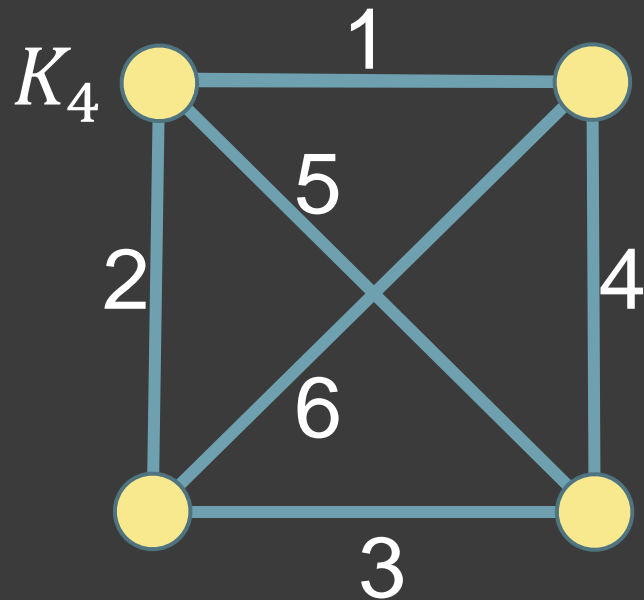
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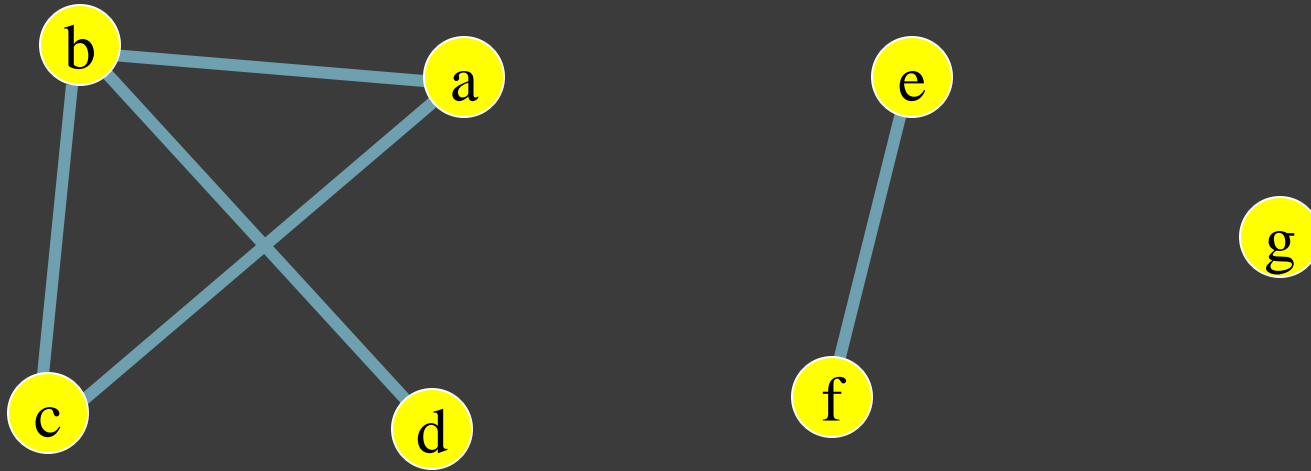
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an instance of
 $G(4, 1/2)$

For $n = 7$,



$$\Pr[G(7, p) = G] = p^5 (1 - p)^{21-5}$$

$$\text{as } \binom{7}{2} = \frac{7 \cdot 6}{2} = 21$$

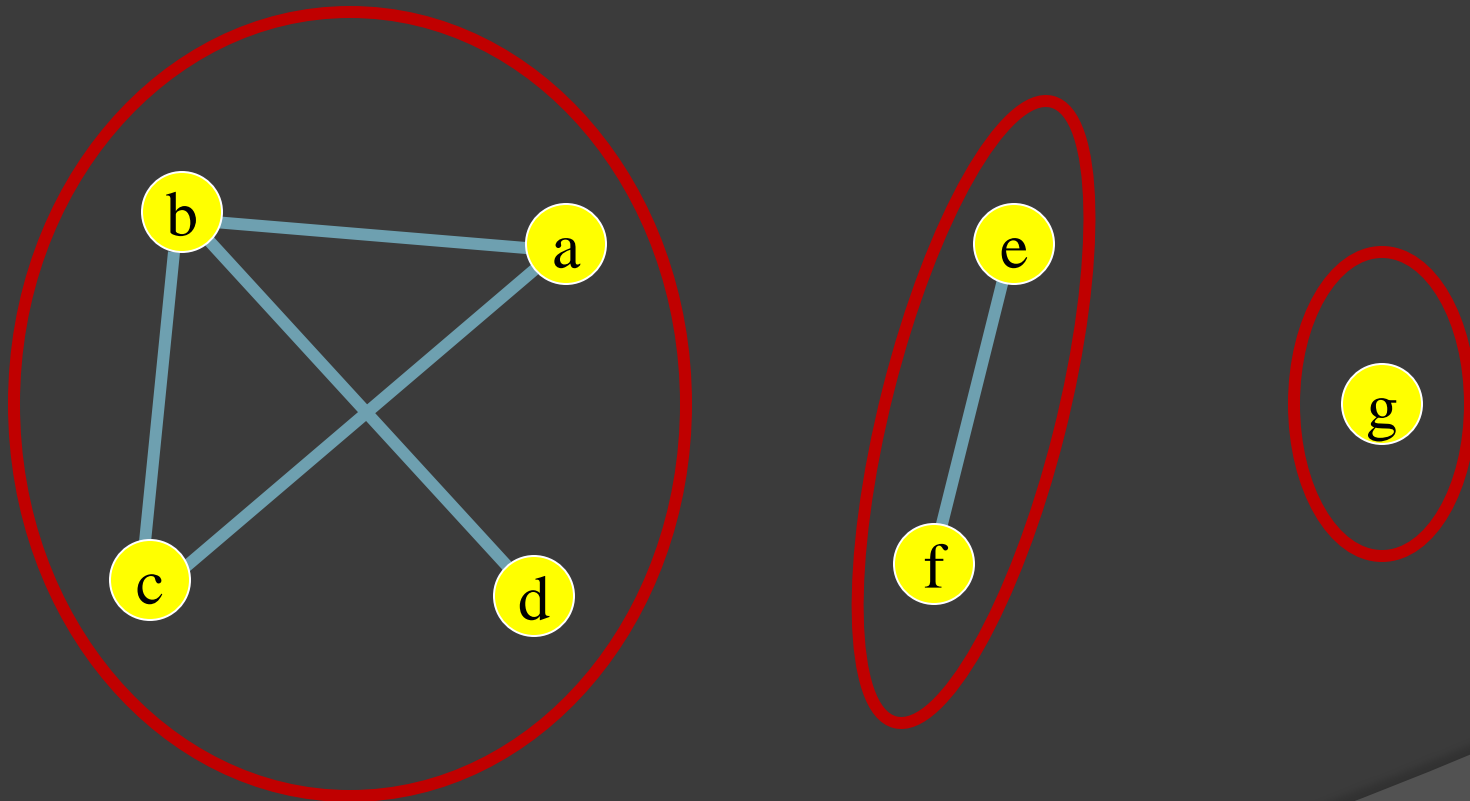
- Expected number of edges in $G(n, p)$:

$$p \binom{n}{2}$$

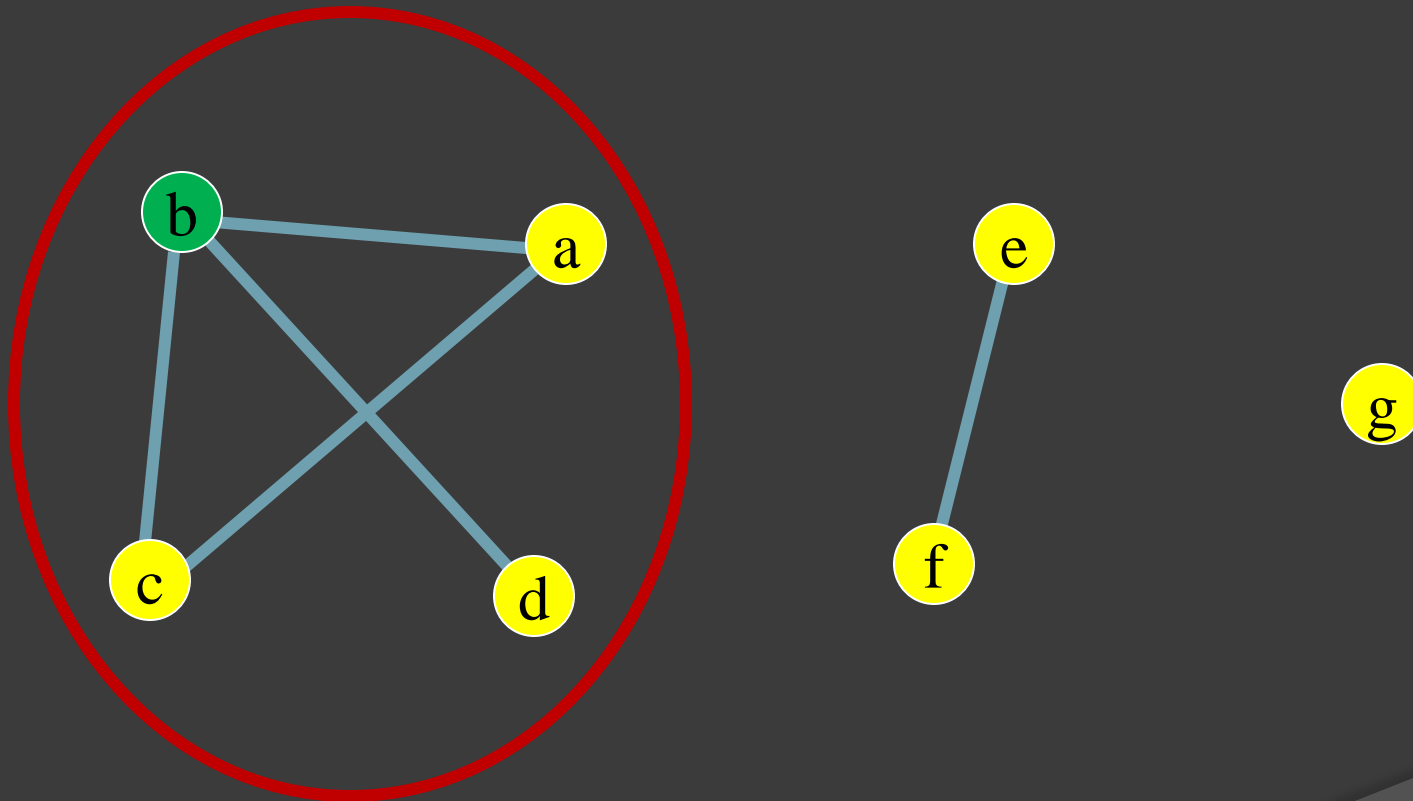
- For a fixed graph G with n vertices and m edges,

$$\Pr[G(n, p) = G] = p^m (1 - p)^{\binom{n}{2} - m}$$

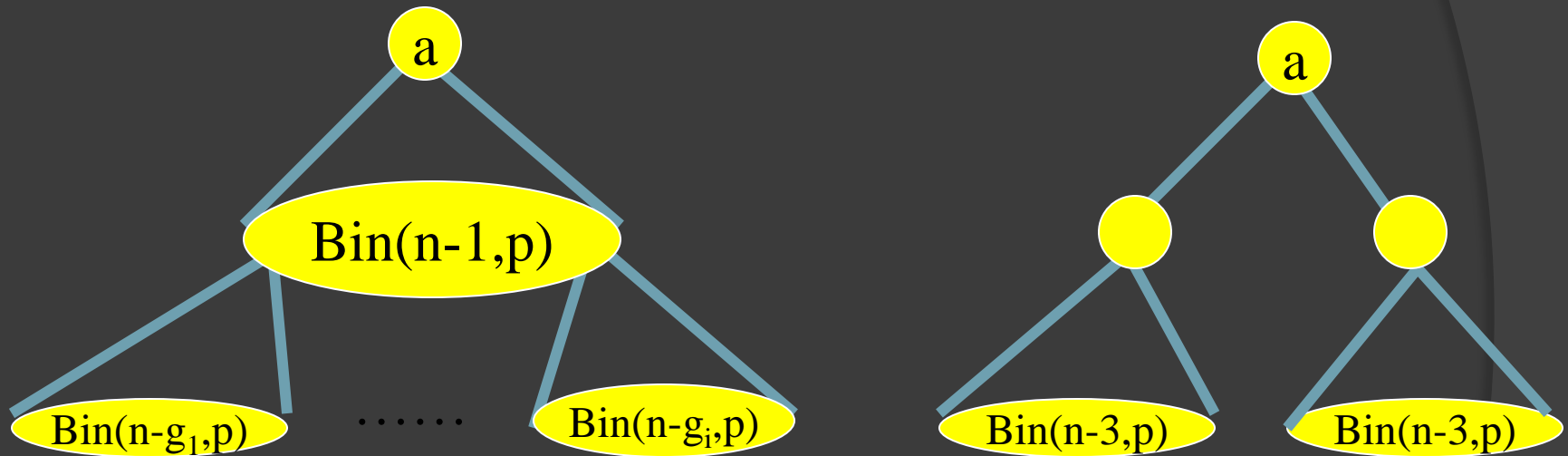
(Connected) Component



Component containing a vertex



Component containing a vertex in $G(n, p)$



where

$$\Pr[\text{Bin}(t, p) = k] = \binom{t}{k} p^k (1 - p)^{t-k}$$

Emergence of Giant Component

- Erdős & Rényi ('60,'61):

In $G(n, p)$ with $pn = \lambda$ for a constant λ

there is a giant component
if and only if $\lambda > 1$

Emergence of Giant Component

- Erdős & Rényi ('60,'61):

For the size $W = W(n, p)$ of a largest component of $G(n, p)$,

$$W = \begin{cases} < c_\lambda \log n & \text{if } pn \rightarrow \lambda < 1 \\ \theta_\lambda n & \text{if } pn \rightarrow \lambda > 1, \end{cases}$$

where θ_λ is the positive solution for

$$1 - \theta - e^{-\theta\lambda} = 0$$

⊙ Łuczak ('90)

In $G(n, p)$ with $pn = \lambda$,

the giant component emerges
when $\lambda - 1 \gg n^{-1/3}$

improving $\lambda - 1 \gg n^{-\frac{1}{3}} \log n$ due to
Bollobás ('84).

• Łuczak ('90)

In $G(n, p)$ with $\lambda = pn$ satisfying

$$\varepsilon := \lambda - 1 \gg n^{-1/3},$$

$$W = (2 + O(\varepsilon)) \varepsilon n$$

and for the size W' of a second largest component

$$W' = \Theta(\varepsilon^{-2} \log(\varepsilon^3 n)) \ll \varepsilon n$$

⦿ Łuczak ('90)

Theorem (supercritical region) Let $\lambda := pn = 1 + \varepsilon$ with $\varepsilon \gg n^{-1/3}$. Then with probability $1 - O((\varepsilon^3 n)^{-1/9})$,

$$|W(n, p) - \theta_\lambda n| \leq 0.2 n^{2/3}.$$

Emergence of Giant Component

- Erdős & Rényi ('60,'61):

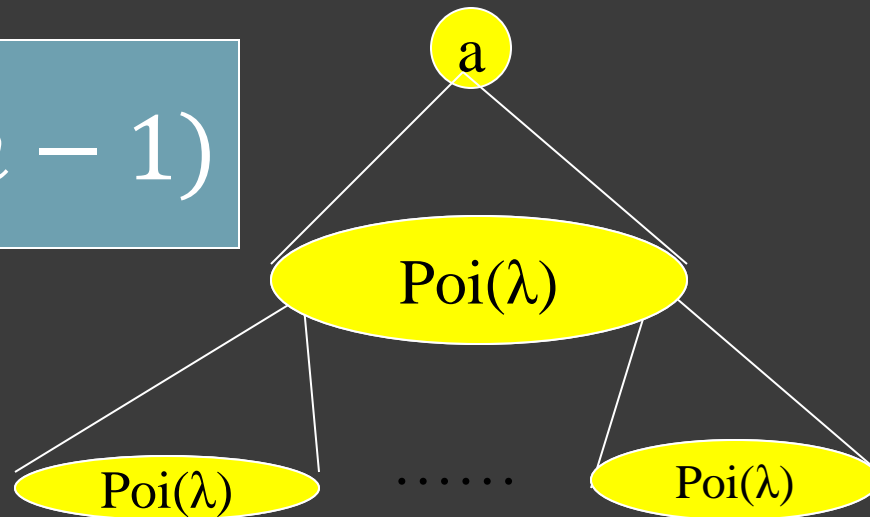
In $G(n, p)$ with $pn = \lambda$ for a constant λ

There is a giant component
if and only if $\lambda > 1$.

Why 1 ?

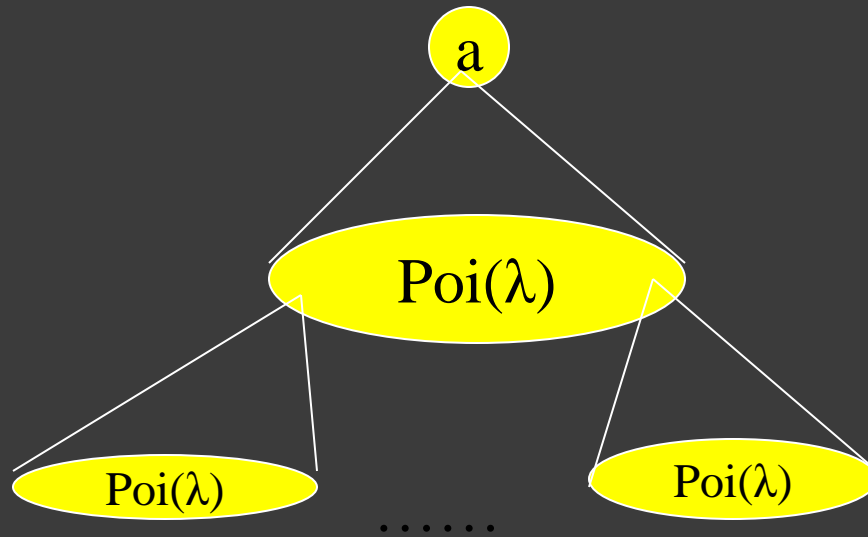
Component containing a vertex in $G(n, p)$

For $\lambda = p(n - 1)$



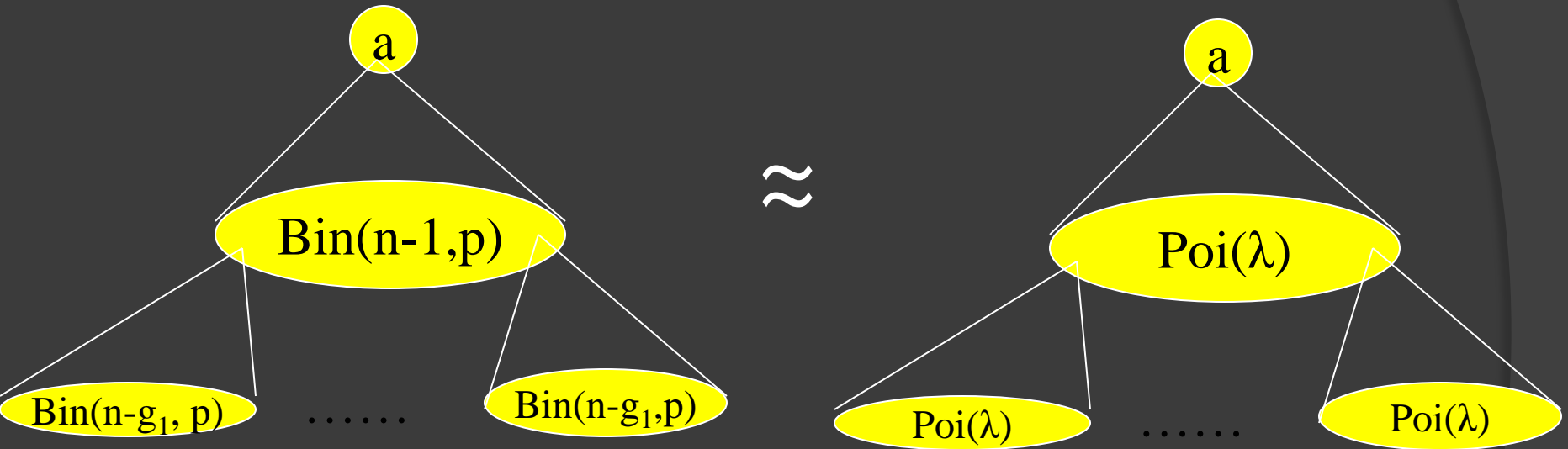
$$\Pr[\text{Poi}(\lambda) = k] = e^{-\lambda} \frac{\lambda^k}{k!}$$

Poisson Branching Process



- If $\lambda < 1$, the process dies out with probability 1
- If $\lambda > 1$, the process survives forever with probability θ_λ (recall θ_λ is the positive solution for $1 - \theta - e^{-\theta\lambda} = 0$).

Two Obstacles



- The difference between $\text{Bin}(n-1, p)$ and $\text{Poi}(\lambda)$.
- The drift : $p(n - g_i)$ keeps decreasing

Similar but more serious obstacles occur in the analyses of

- The k -core problem of the random graph (Pittel-Spencer-Wormald, ...)
- Giant strong component of the random directed graph (Karp,...)
- Pure literal algorithm for the random satisfiability problem (Broder-Frieze-Upfal, ...)
- Unit clause algorithm for the random satisfiability problem (Chao-Franco,...)
- Karp-Sipser Algorithm to find a large matching in the random graph (Karp-Sipser, ...)

The First Obstacle

- The difference between

$\text{Bin}(n - 1, p)$ and $\text{Poi}(\lambda)$.

More generally, is there a random graph model $G_{\text{new}}(n, p)$, in which all degrees are i.i.d $\text{Poi}(\lambda)$ and

$$G_{\text{new}}(n, p) \approx G(n, p)?$$

NO: The sum of degrees must be even.

YES in an asymptotic sense

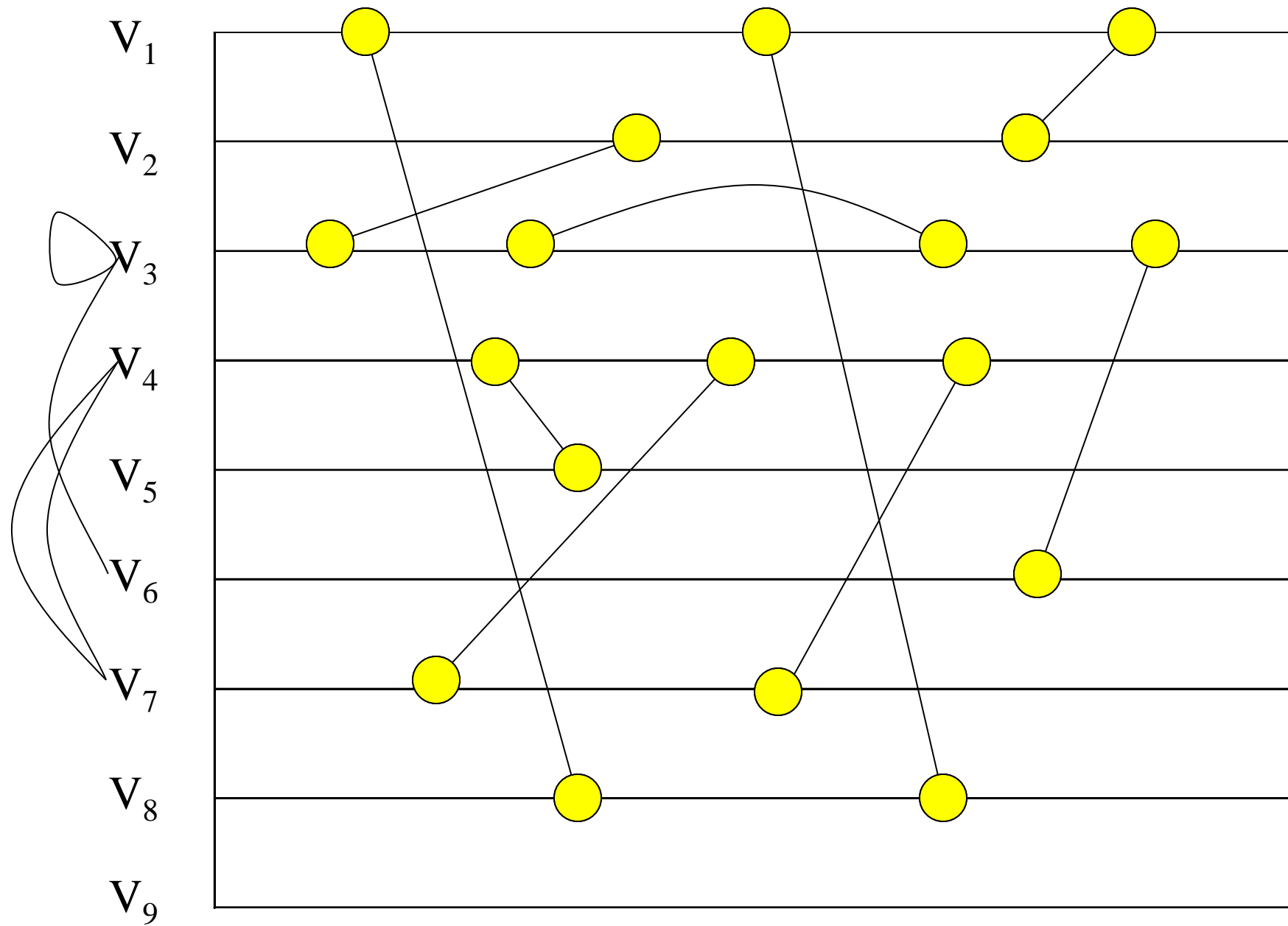
Is it possible to define
a random graph model
in which all degrees
are i.i.d $Poi(\lambda)$?

Just No! It!

Poisson Cloning Model $G_{PC}(n, p)$:

Definition

- Take i.i.d Poisson random variables $d(v)$'s, v in V , with mean $\lambda = p(n - 1)$.
- For each vertex v in V , take $d(v)$ copies, or **clones**, of v .
- If $\sum d(v)$ is even, generate a uniform random perfect matching on all clones and then contract clones of the same vertex.



- If $\sum d(v)$ is odd, generate a perfect matching excluding a clone. The excluded clone induces a loop.

(When $d(v) = d$ for all v , this is the configuration model for random regular graphs due to B. Bolloás ('84).)

$$G(n, p) \approx G_{PC}(n, p)$$

Theorem (K) If $pn = O(1)$, then, for any event A regarding $G(n, p)$,

$$Pr[A] \geq c_1 Pr_{PC}[A] - e^{-(pn^2)}$$

and

$$Pr[A] \leq c_2 Pr_{PC}[A]^{1/2} + e^{-\Omega(pn^2)}$$

In particular,

$$Pr[A] \rightarrow 0 \quad \text{iff} \quad Pr_{PC}[A] \rightarrow 0 \quad \text{and}$$

$$Pr[A] \rightarrow 1 \quad \text{iff} \quad Pr_{PC}[A] \rightarrow 1$$

The Second Obstacle

- ⦿ The drift : $p(n - g_i)$ keeps decreasing

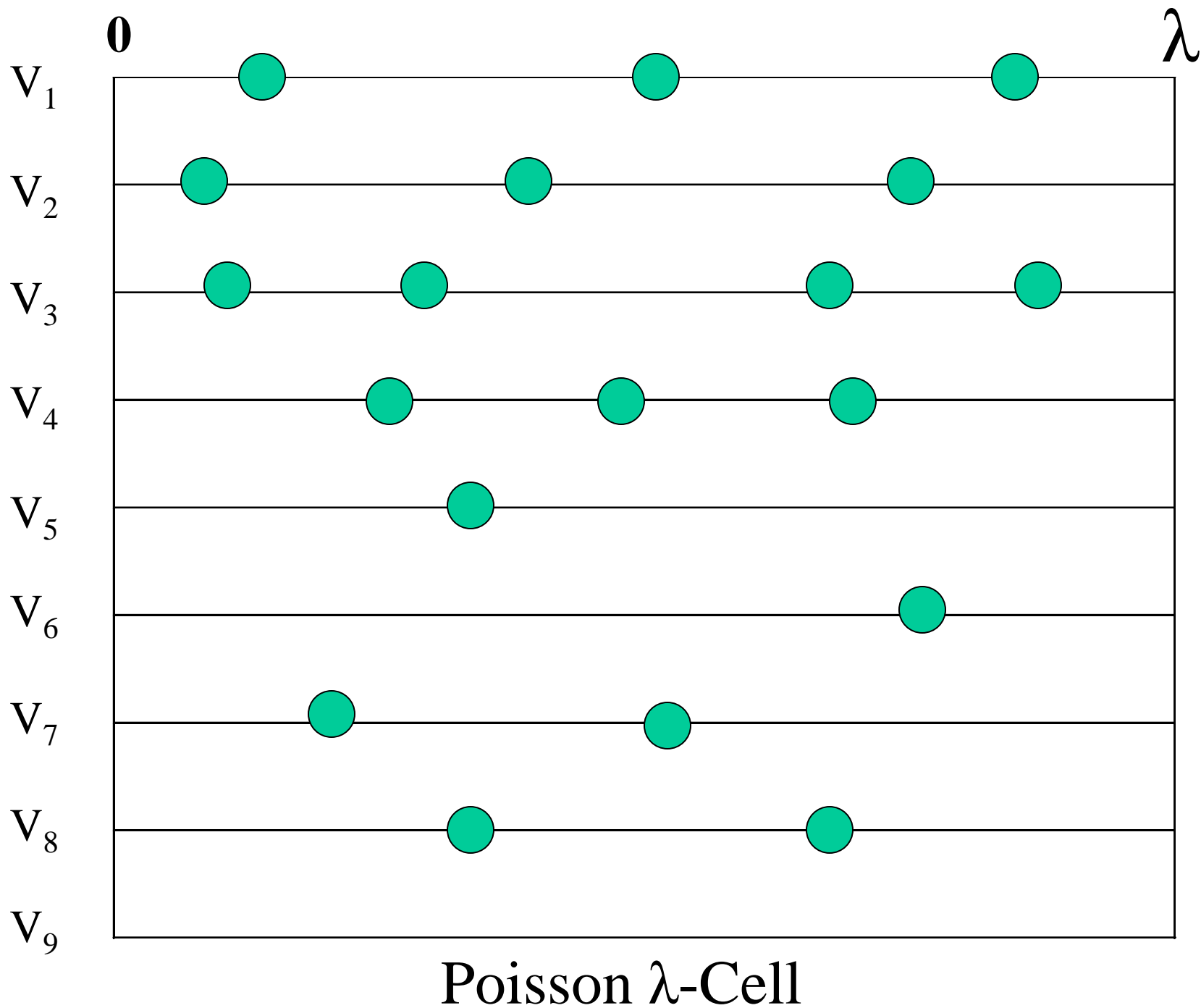
We introduce

the cut-off line algorithm

(COLA)

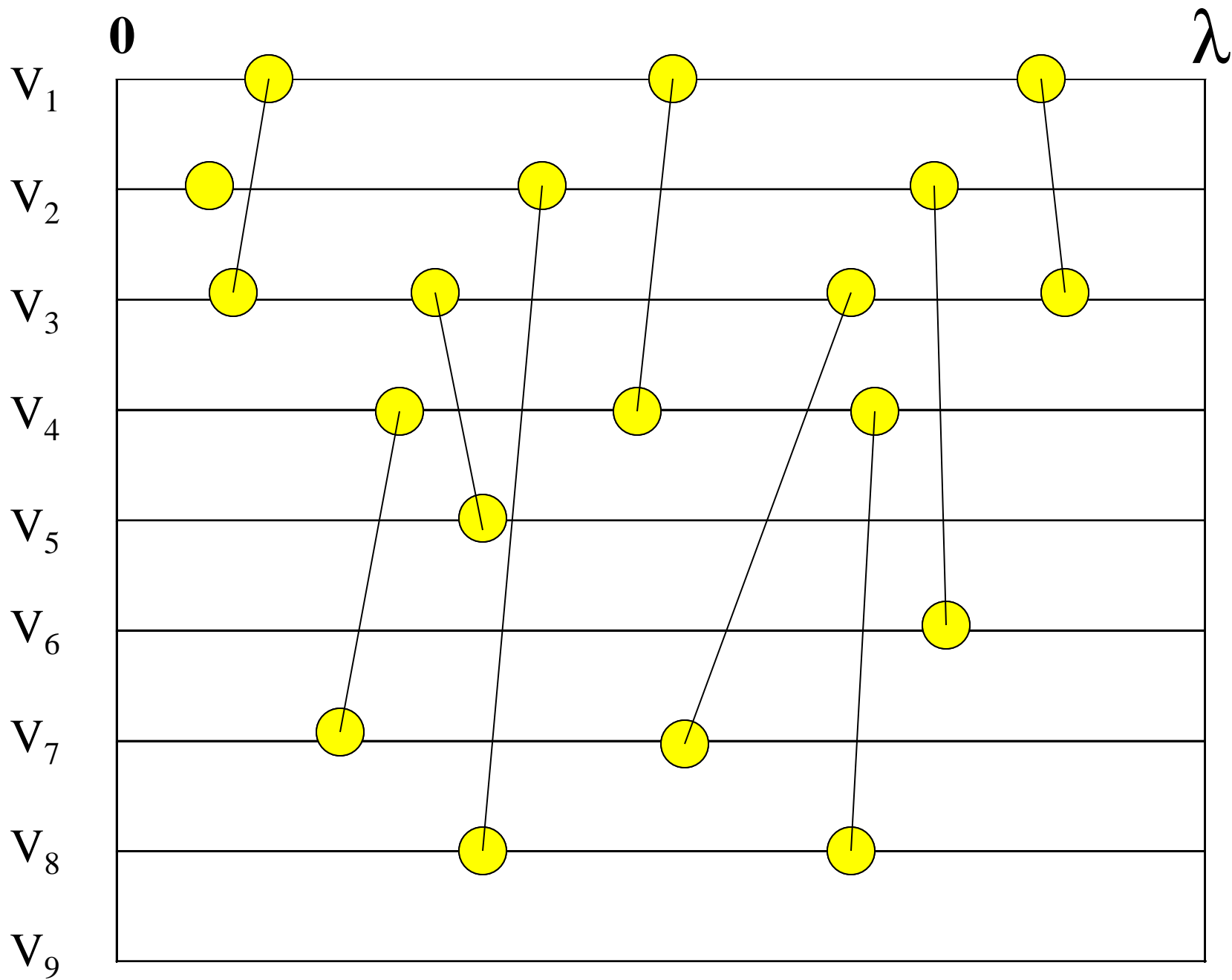
Poisson λ -cell

- ⦿ For each clone w , assign a uniform random (real) number between 0 and λ , independently of all others.
- ⦿ A clone is larger than another clone if so are the assigned numbers.



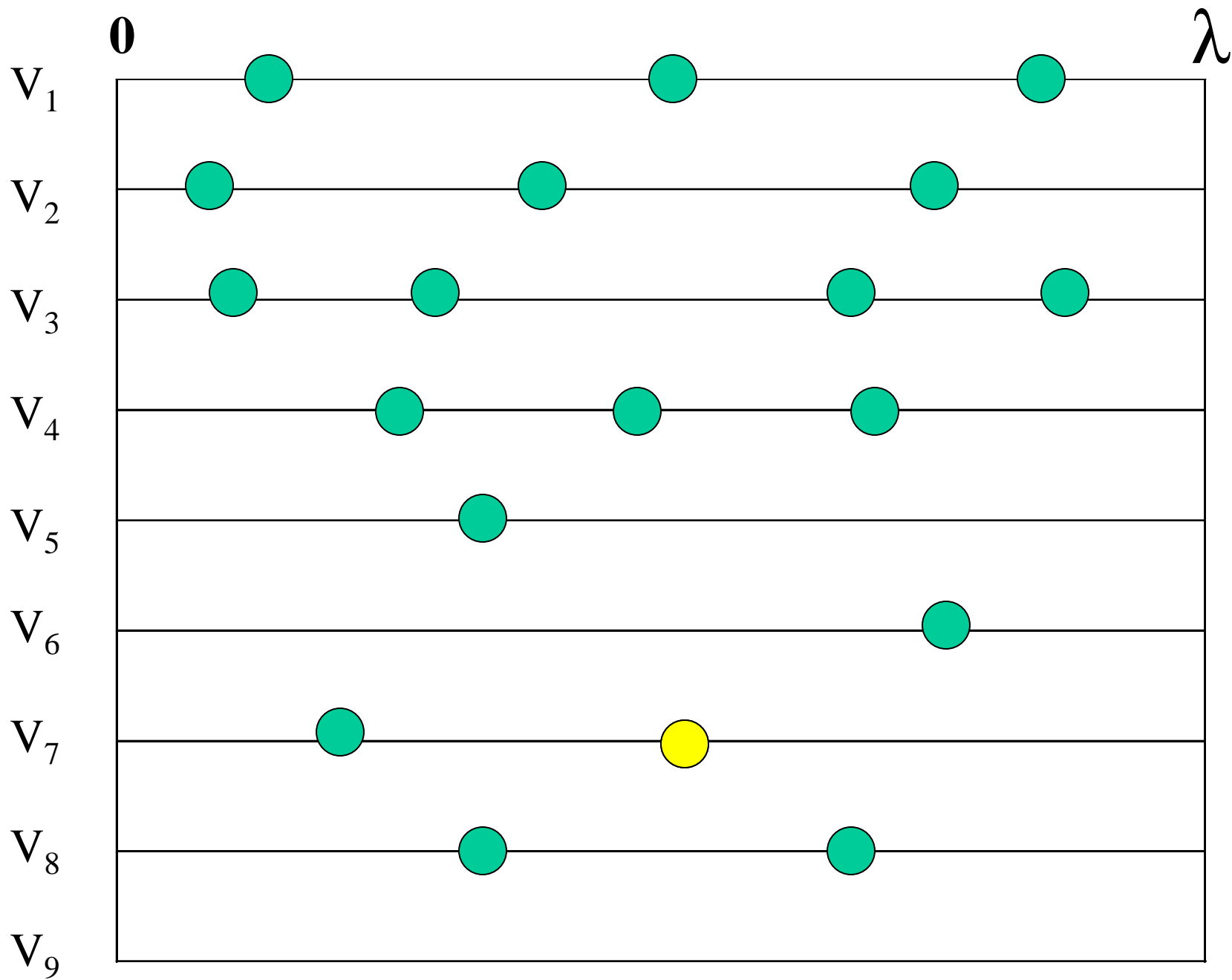
A way to generate
the random perfect matching

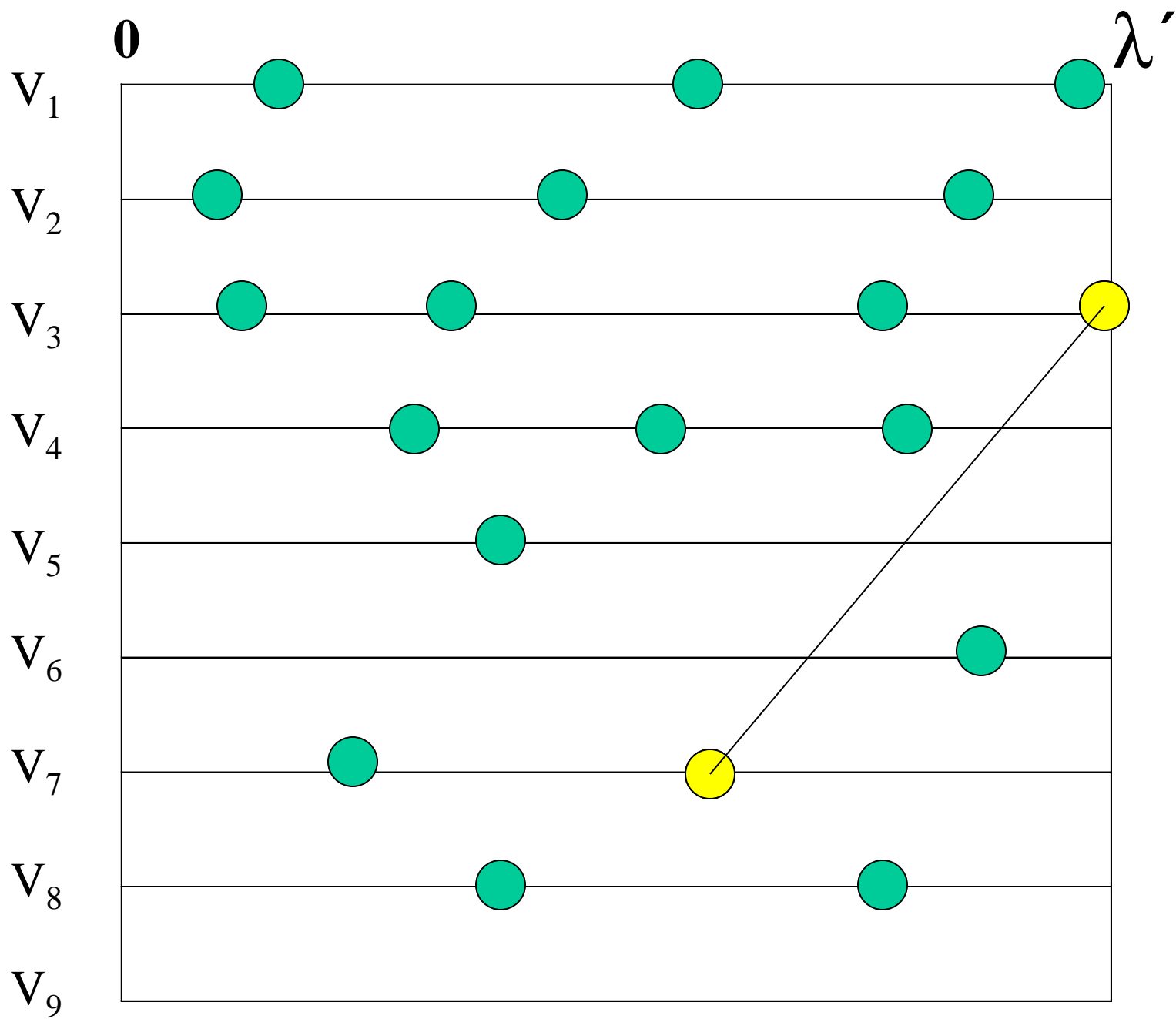
Take two largest clones and match them.
Repeat this for the remaining
unmatched clones



Cut-Off Line Algorithm (COLA)

Choose the first unmatched clone
according to a certain selection rule
independently of assigned numbers
and
match it to the largest
unmatched clone



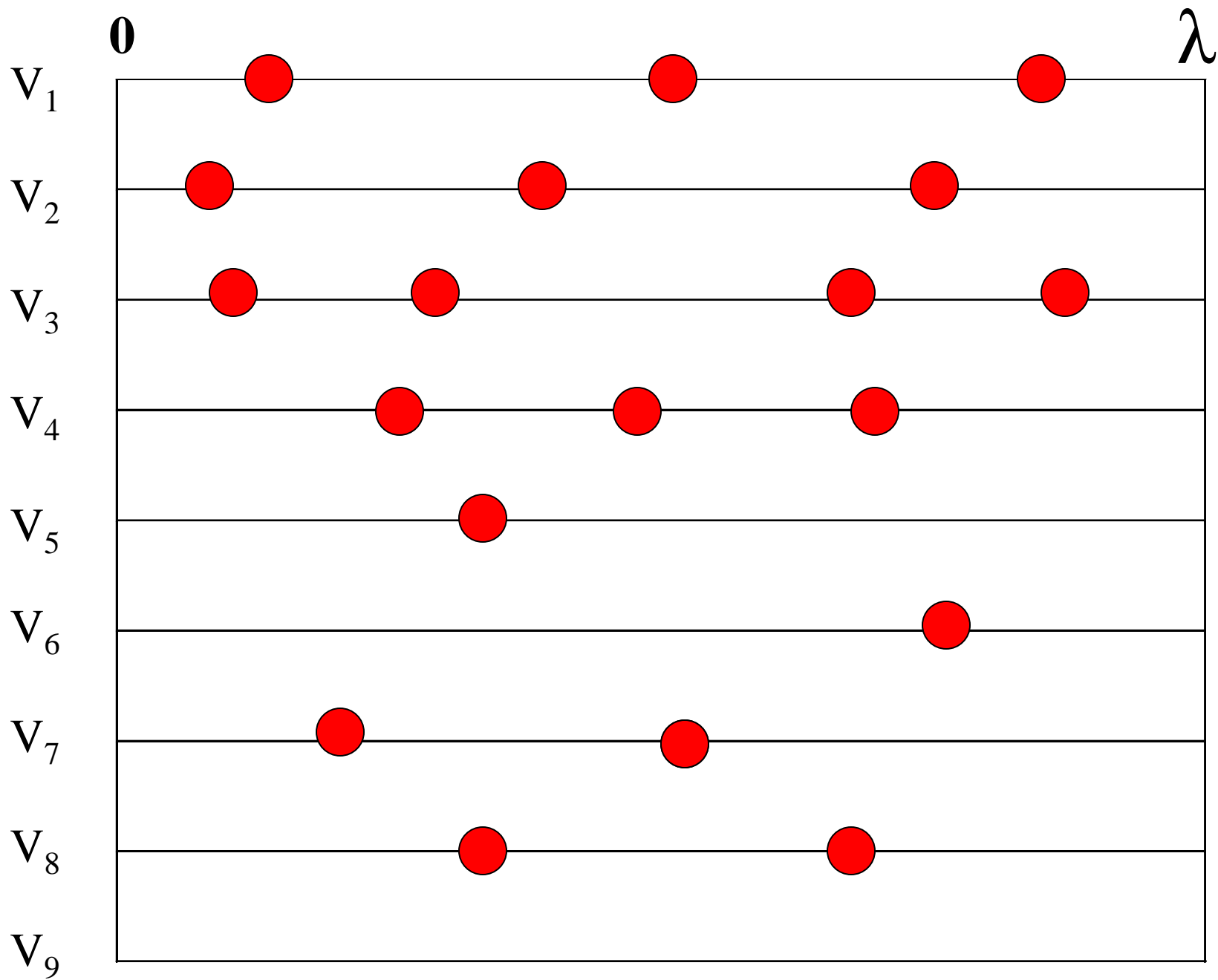


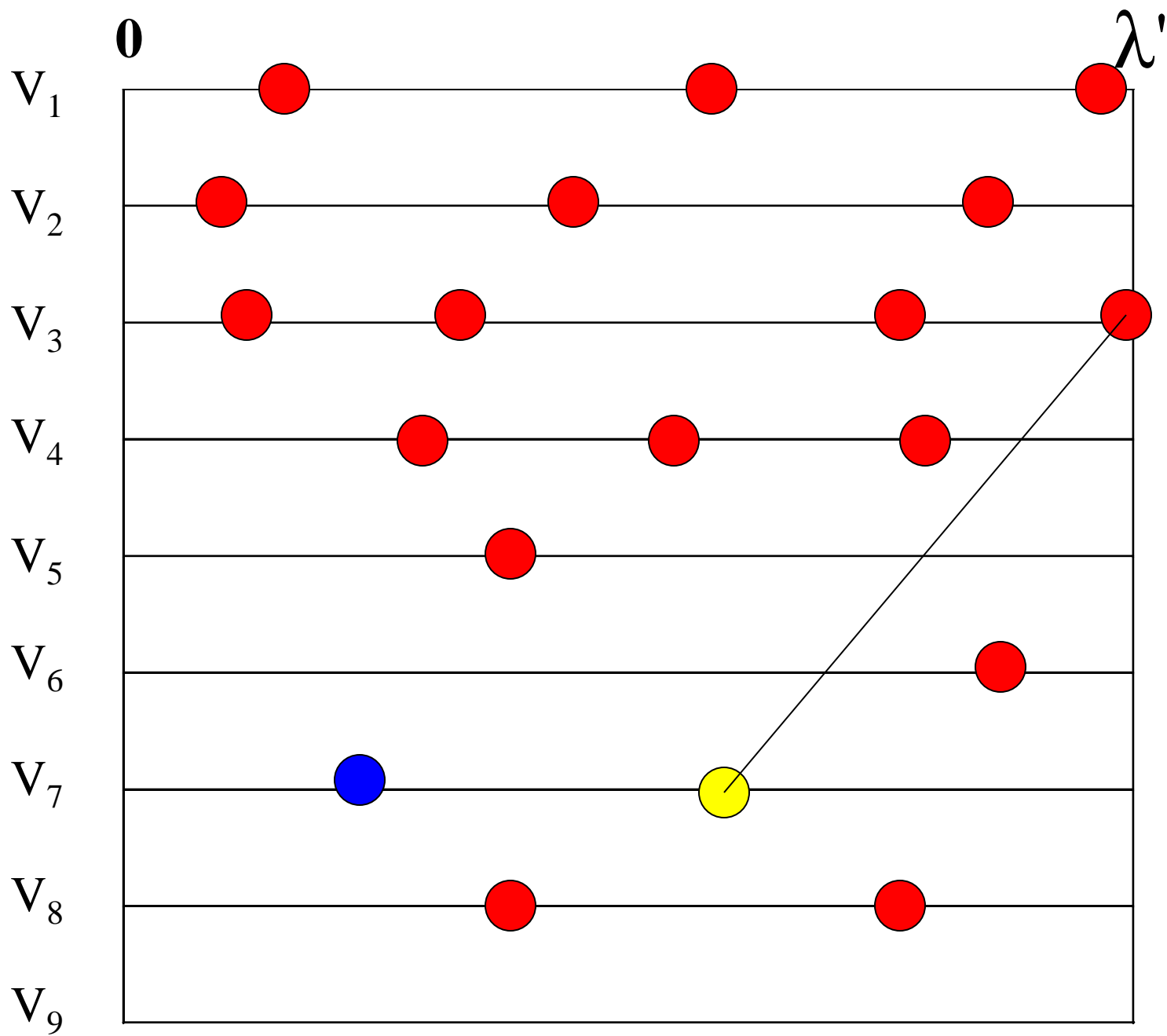
Cut-Off Line Algorithm (COLA)

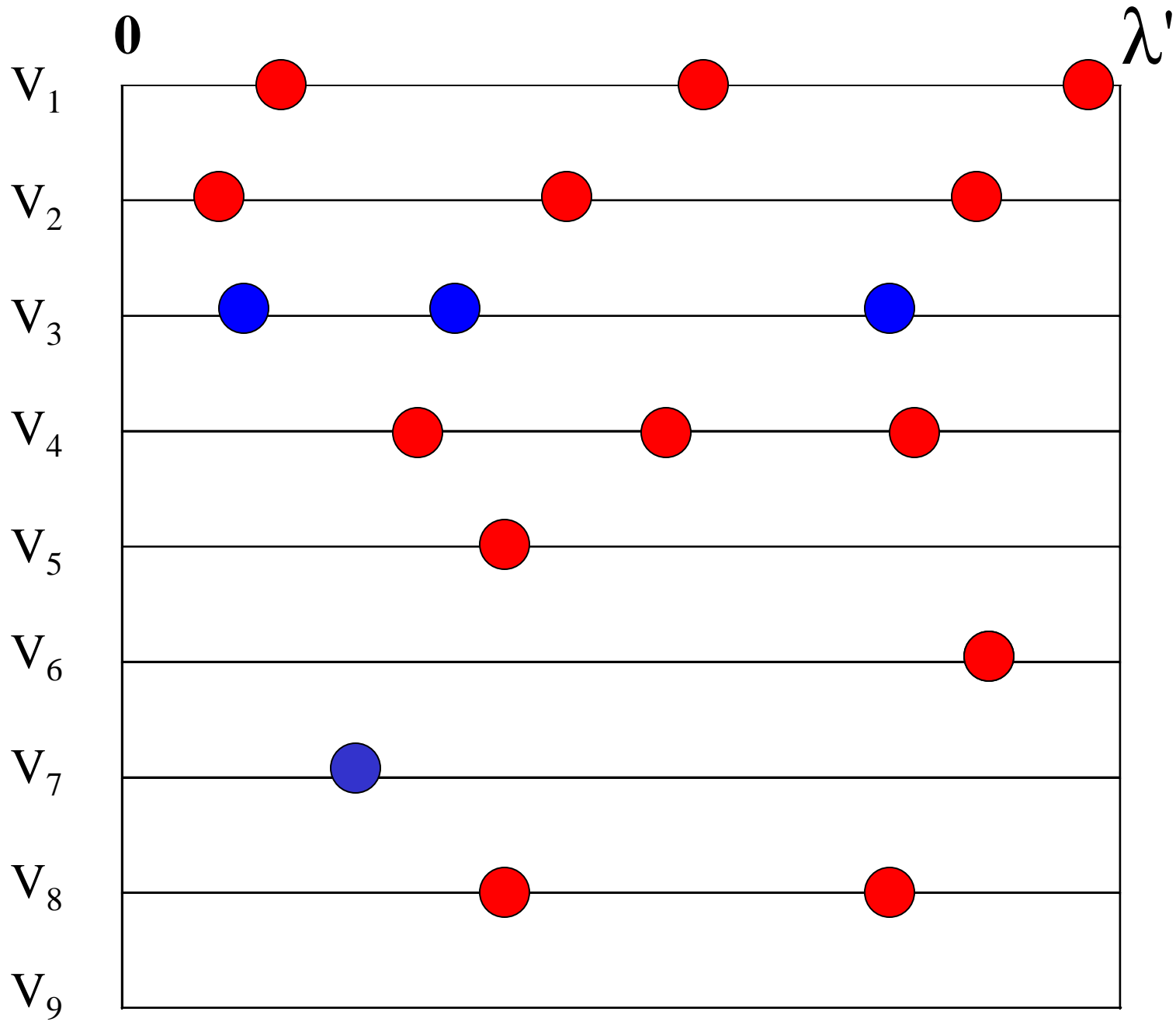
- Initially, the cut-off value $\Lambda = \lambda$. The cut-off line is the vertical line containing $(\Lambda, 0)$.
- After one step, the cut-off value $\Lambda = \lambda'$. The cut-off line is the vertical line containing $(\Lambda, 0)$.

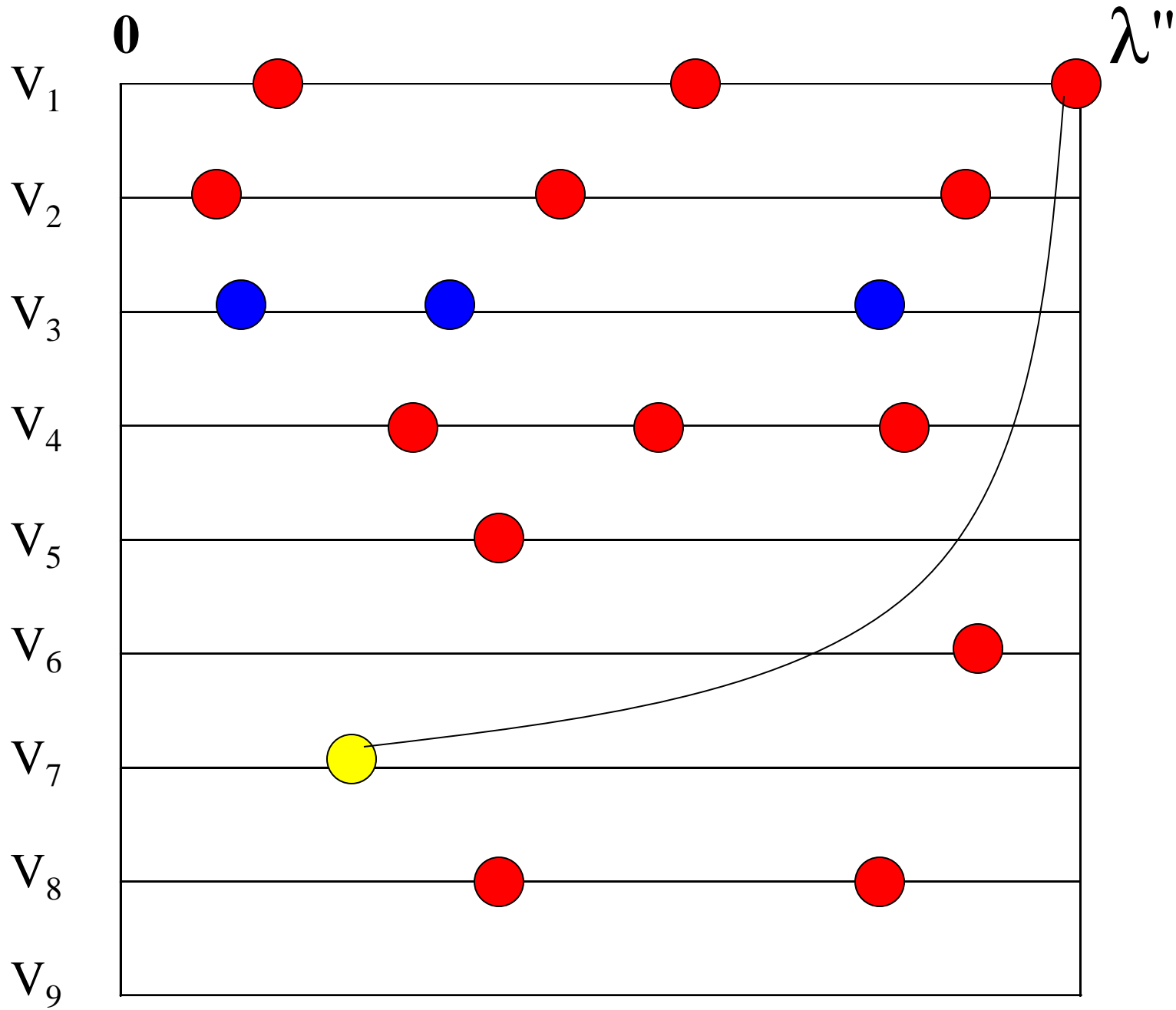
Notice that λ' is the largest number among $N - 1$ independent uniform random numbers between 0 and λ .

Cut-Off Line Algorithm for Component









New Results

Random Graph $G(n, m)$

- For a fixed vertex set V of n elements, consider all graphs on V with m edges.

There are

$$\binom{\binom{n}{2}}{m} \text{ such graphs}$$

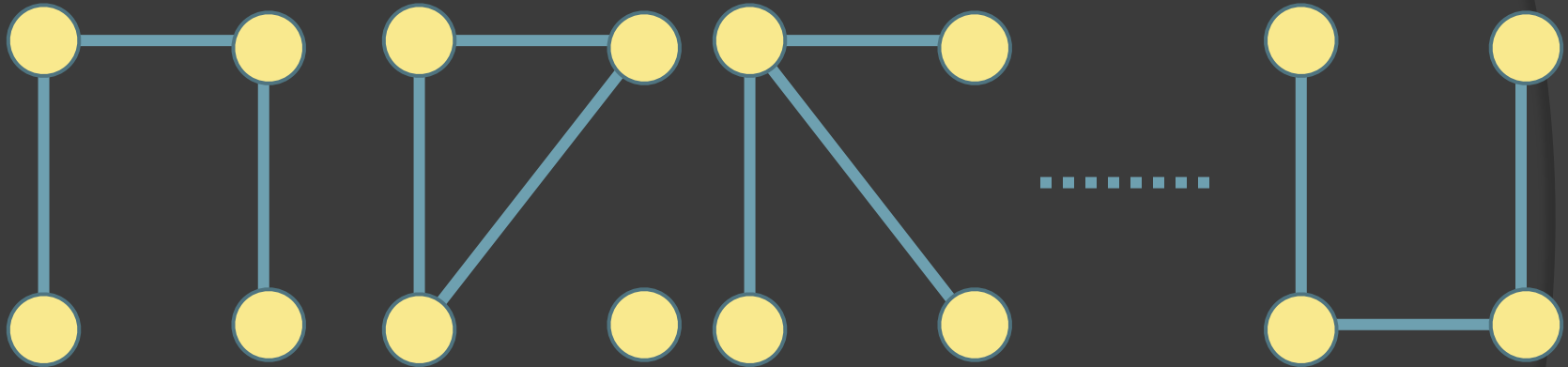
- Choose a graph uniformly at random from all such graphs.

- Easy

$$G(n, m) \approx G(n, p)$$

provided $p \binom{n}{2} = m$.

- For example, if $n = 4$ $m = 3$, then
each of $\binom{\binom{4}{2}}{3} = \binom{6}{3} = 20$ graphs



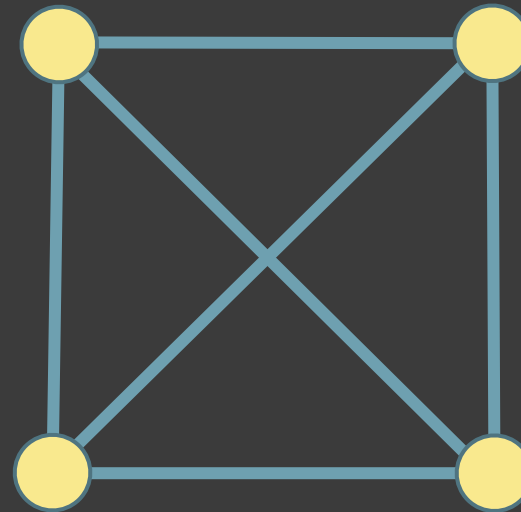
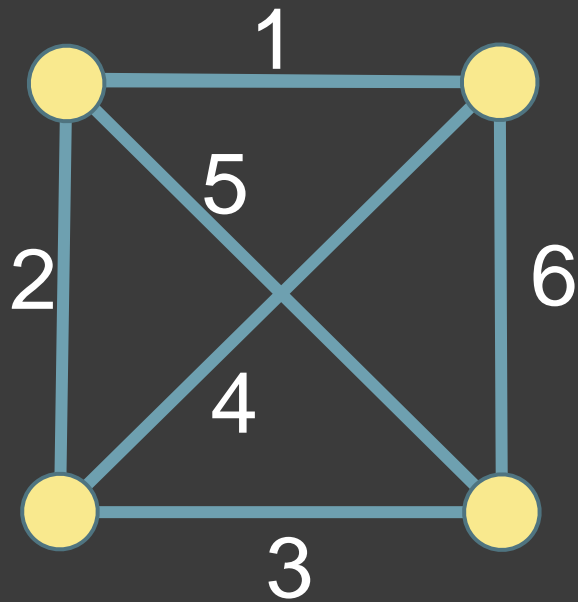
is equally likely to be $G(4,3)$.

Graph Process

$$G(n, 1), \dots, G(n, m), \dots$$

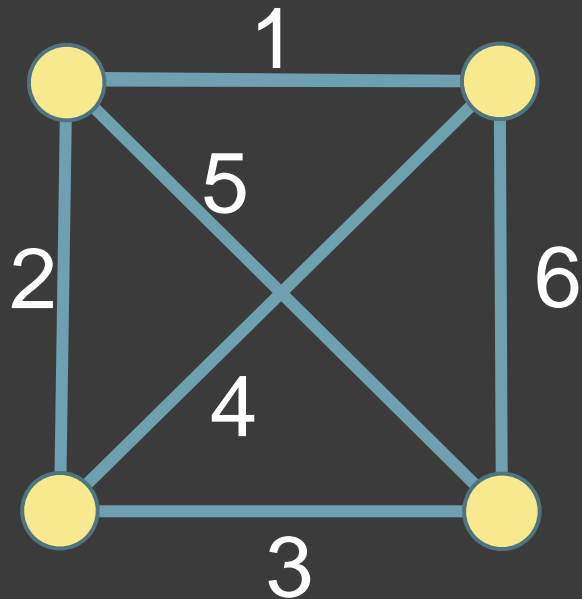
- Randomly order all $\binom{n}{2}$ edges of K_n so that each of $\binom{n}{2}!$ orderings is equally likely.
- The graph $G(n, m)$ is the graph consisting of the first m edges in the random ordering.
- Notice that $G(n, m)$ in the graph process has the same distribution as $G(n, m)$ defined earlier.

G_1 G_2 G_3 G_4 G_5 G_6

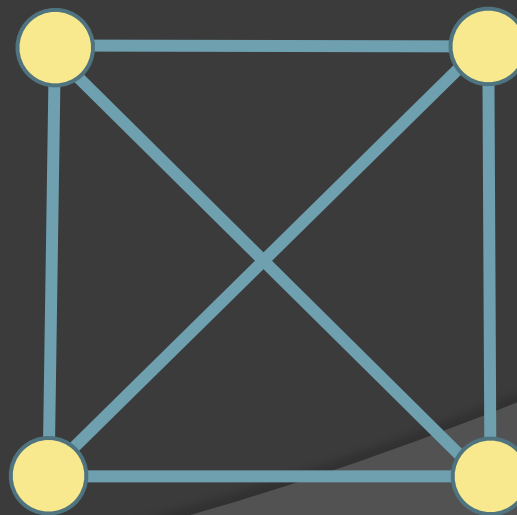


Graph process under constraint

Triangle Free Process: no triangle is allowed



G_1 G_2 G_3 G_4

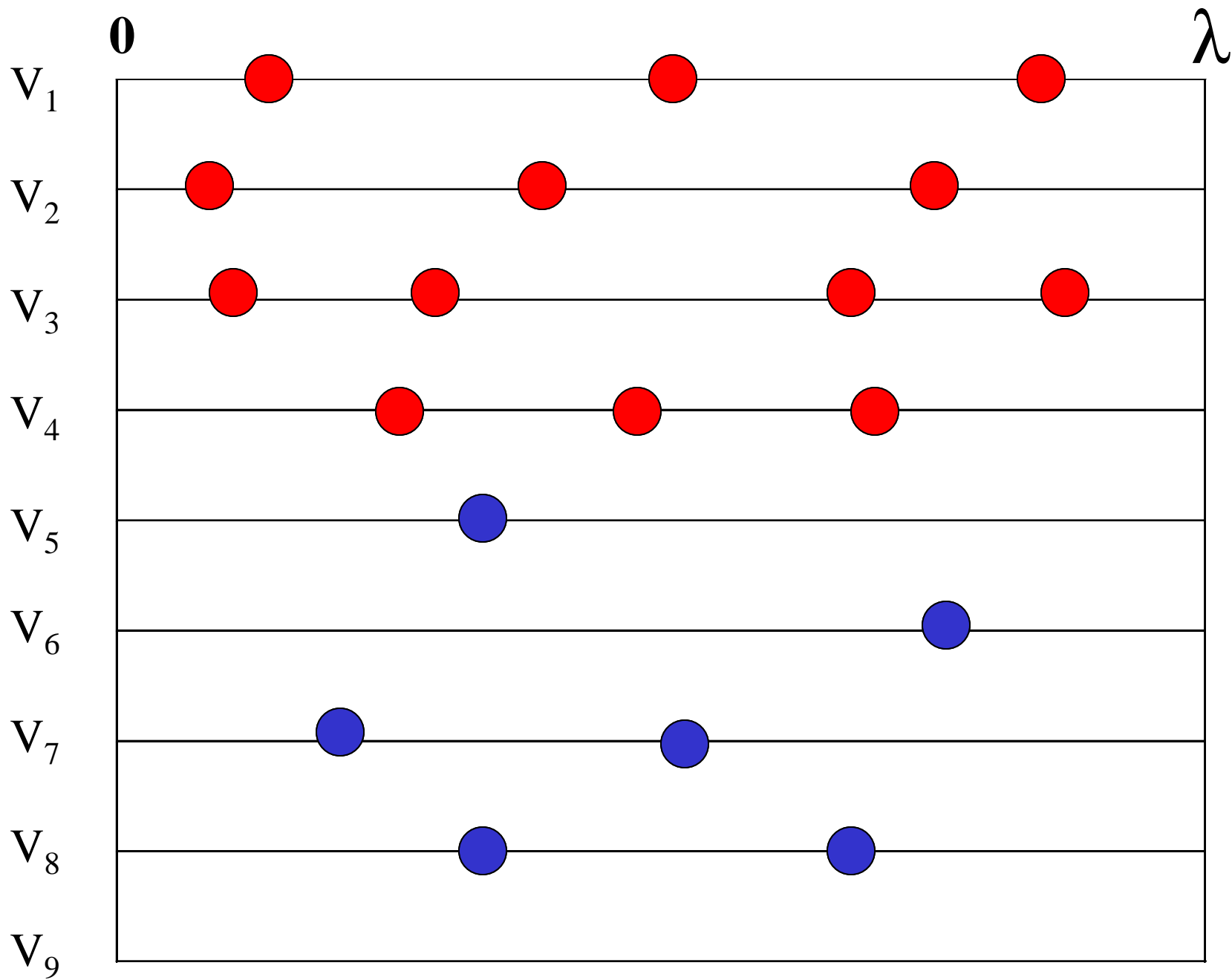


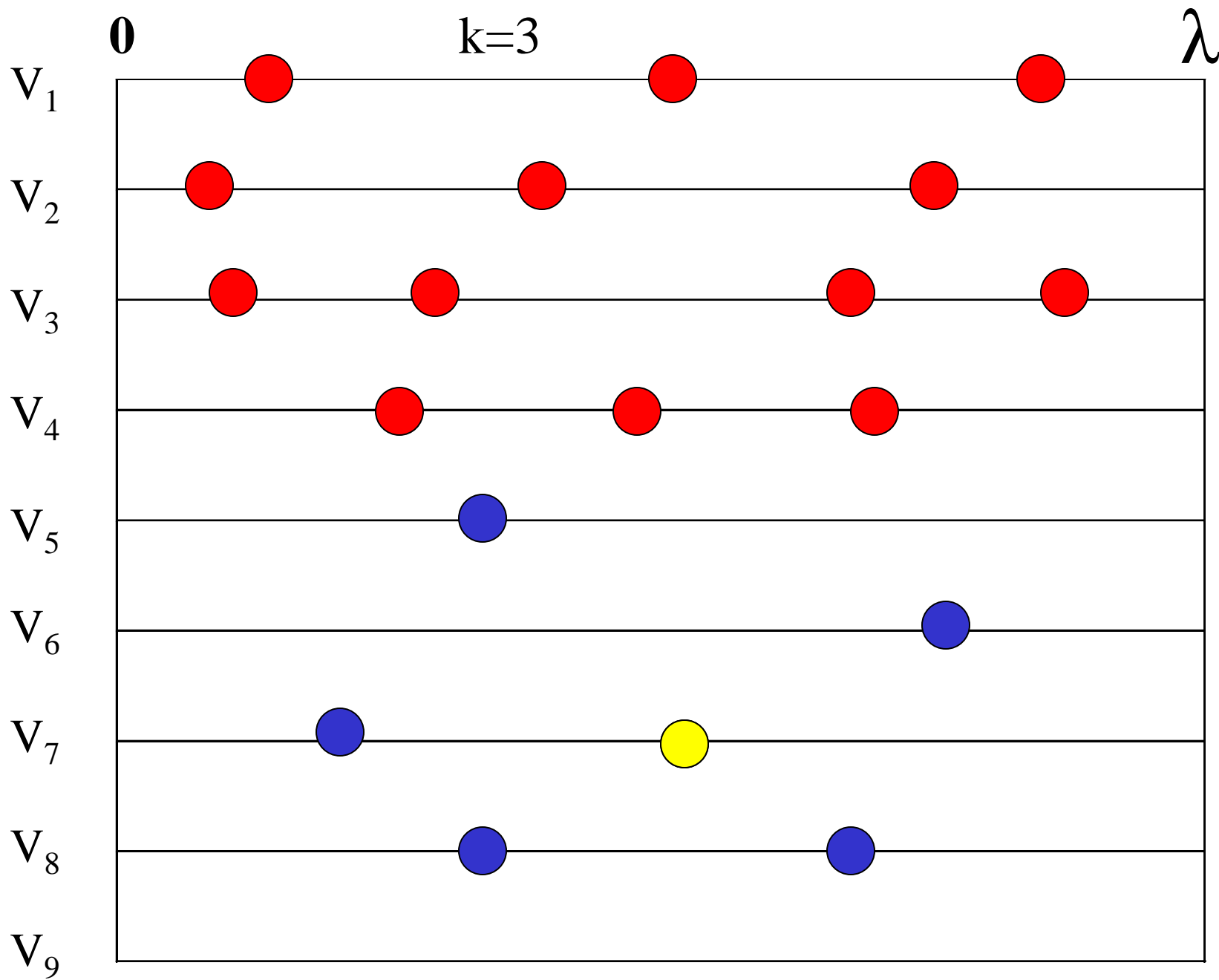
This is called the triangle-free process.

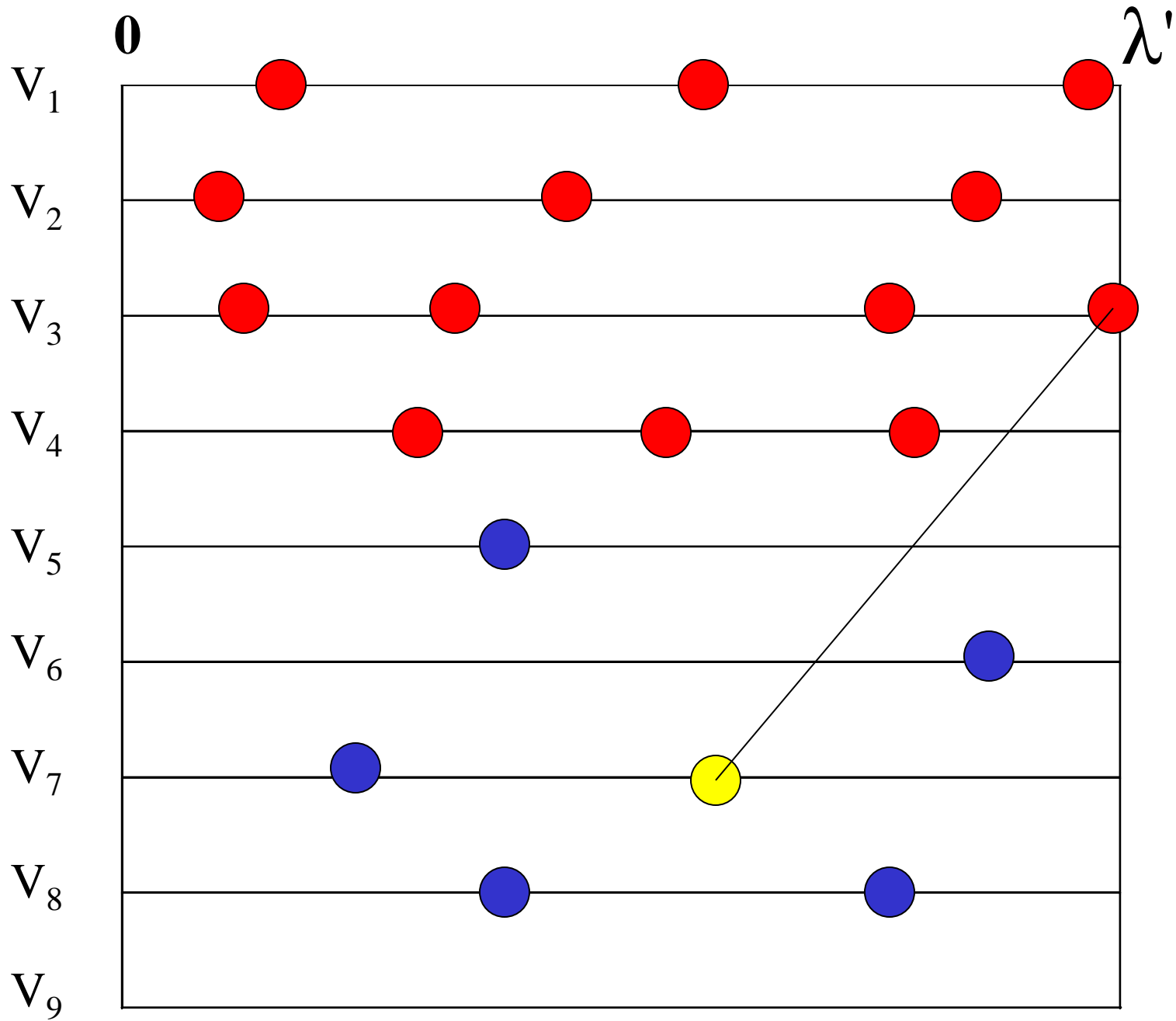
Graph process under constraint

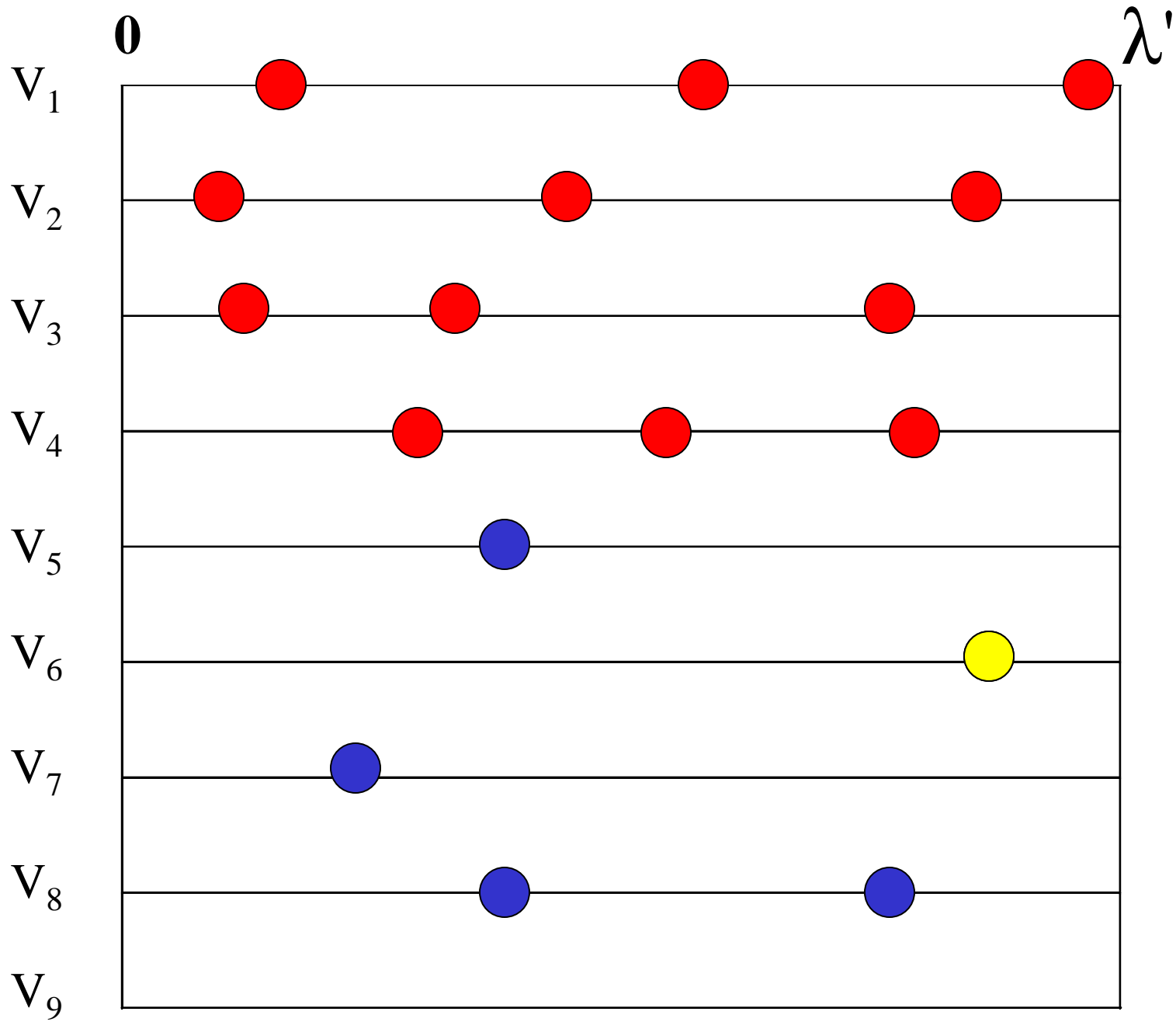
- Many applications to construct graphs satisfying certain properties.
- It is generally not easy to analyze the process
- Studied by many researchers including Ajtai, Komlós, Szemerédi, Rodl, Kahn, and more.

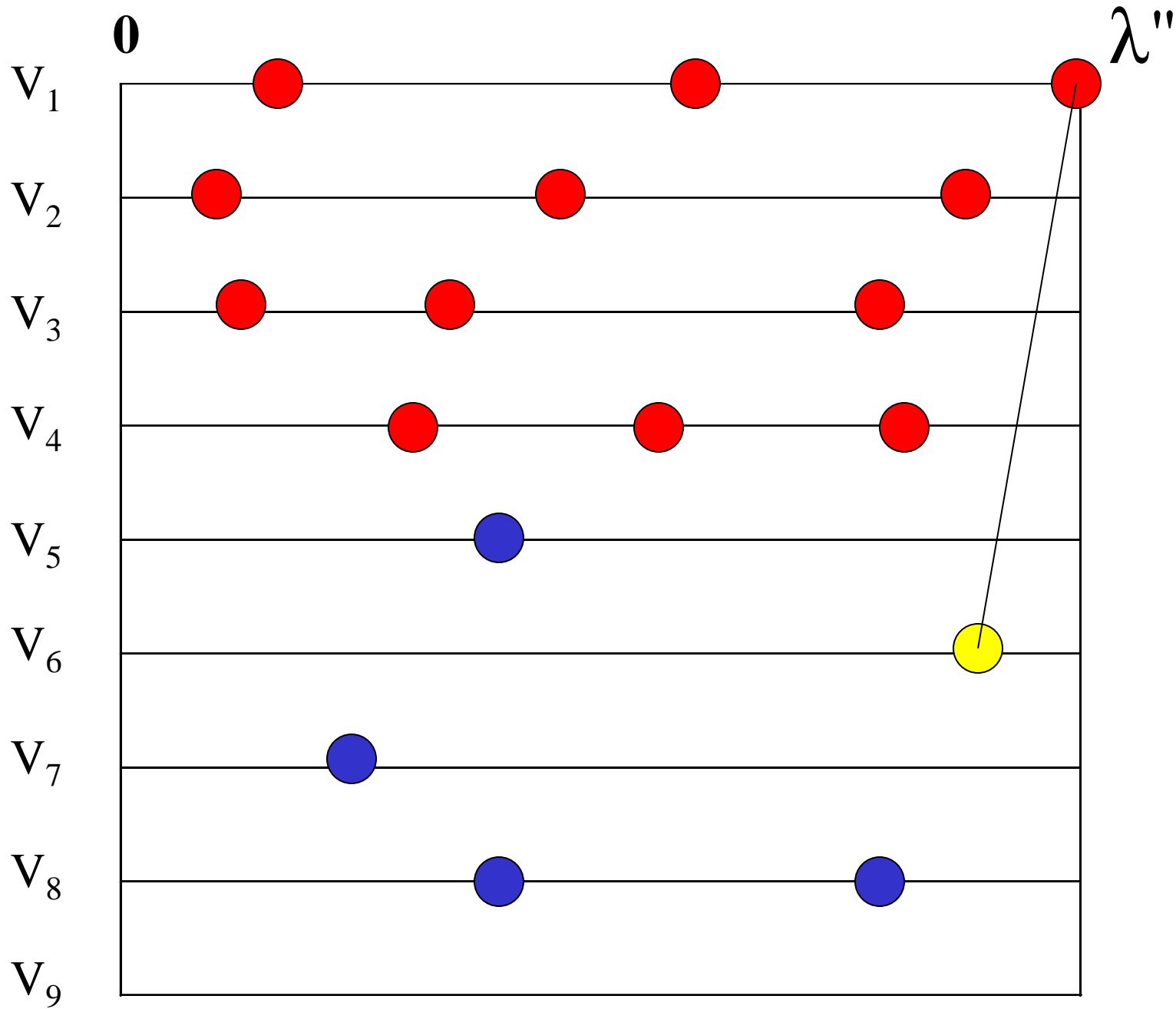
Application: Ramsey number $R(3, t)$

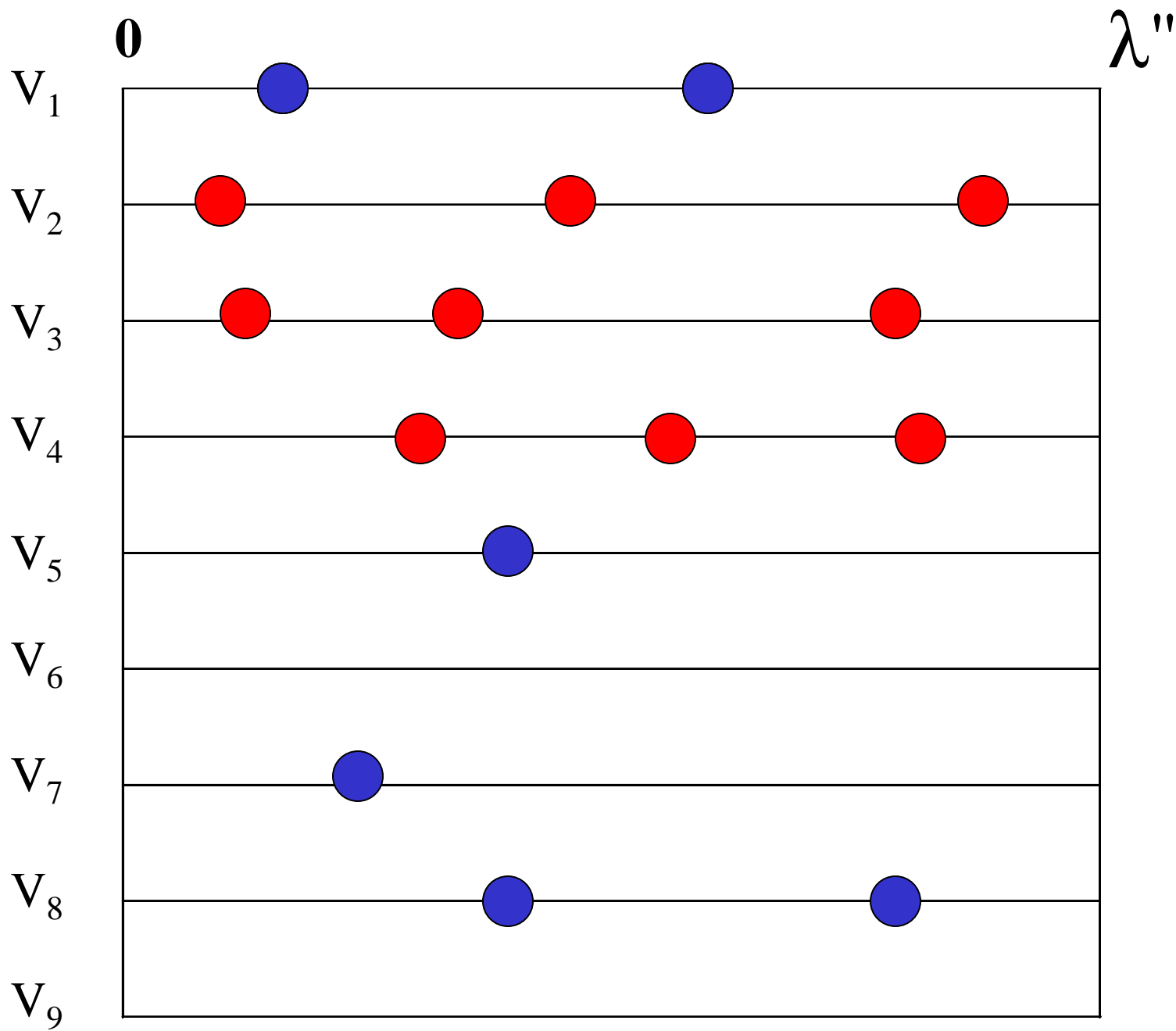


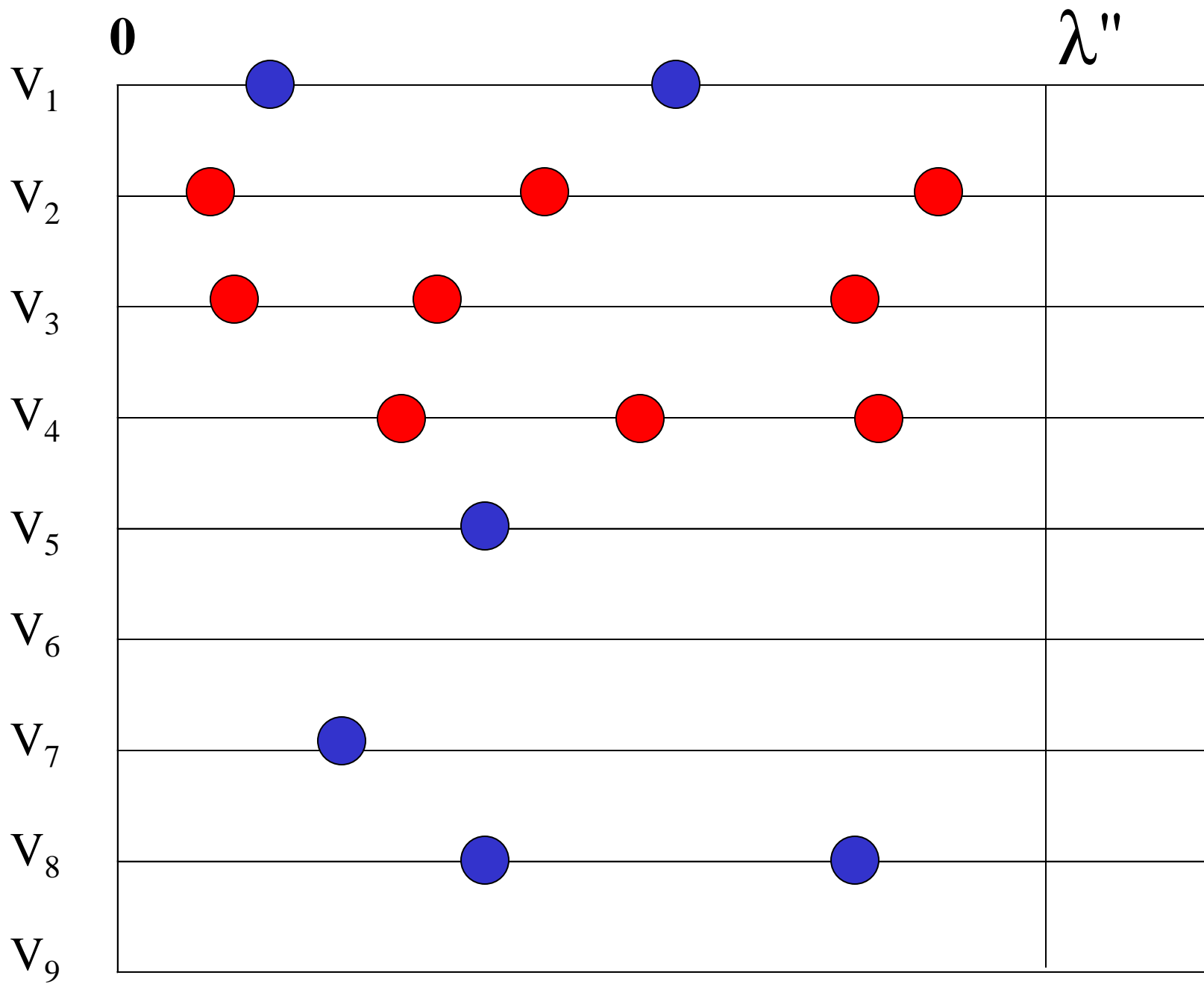






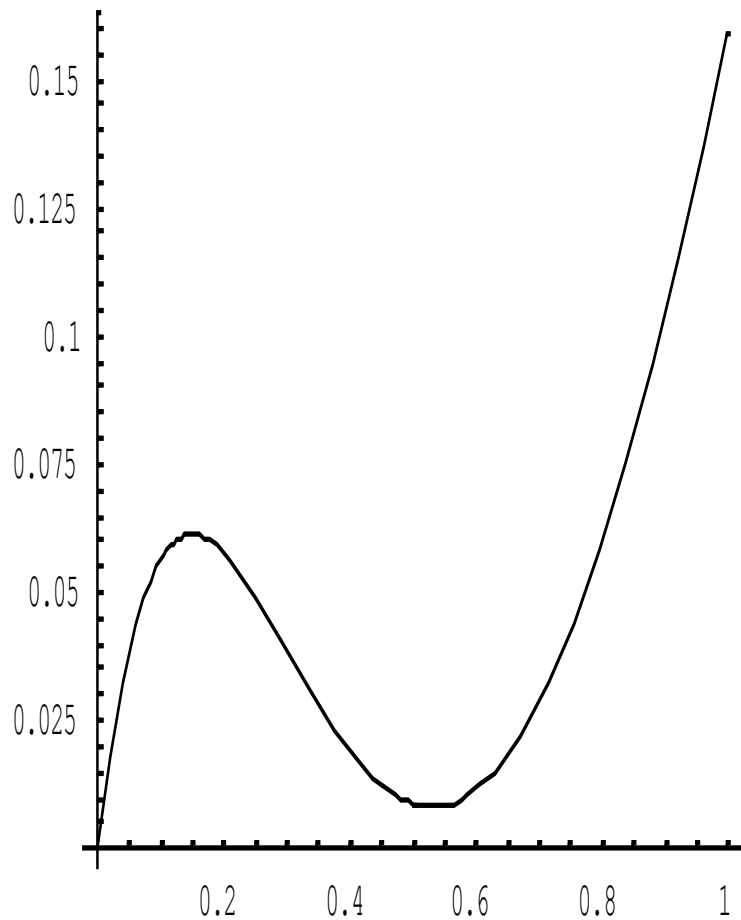




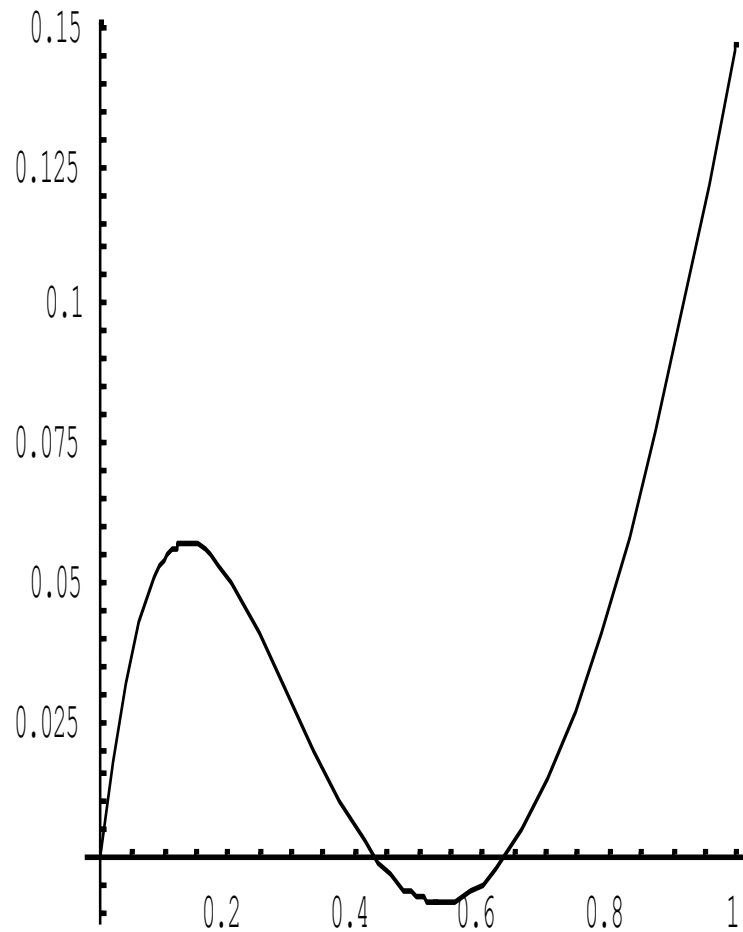


Number of light clones

$k=3, \lambda=3.3$



$k=3, \lambda=3.4$



$$k=3, \lambda=3.35$$

